

# *Utilization and Evolution of KLEE-based Technologies for Embedded Software Testing at Fujitsu*

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## ■ Director of Research

- Fujitsu Labs of America, Sunnyvale, California
- Lead the “Software Quality and Security Lab” – 10 researchers
- Have been with Fujitsu 20 years
  - In Silicon Valley time that is close to 200

## ■ PhD, Princeton University, EE

## ■ Worked in

- Hardware Test (PhD)
- Hardware Verification/Validation (early 2000s)
- Software Validation and QA (from 2006)

# The Fujitsu Group



**FUJITSU** (Est. June 1935)

(70% software related products & services)

Revenue: \$41 billion

R&D Expenditure: \$2.5 billion

## Fujitsu Laboratories Group

- Capital: 5 billion JPY
- R&D Budget: Approx. 30 billion JPY
- Employees: Approx. 1,500 worldwide

Fujitsu Laboratories Ltd.  
Kawasaki Laboratories  
(Japan)  
(Established 1968)

Fujitsu Laboratories Ltd.  
Atsugi Laboratories  
(Japan)  
(Established 1983)

Fujitsu Research and  
Development Center  
Co., Ltd. (China)  
(Established 1998)

Fujitsu Laboratories  
of America, Inc.  
(U.S.)  
(Established 1993)

Fujitsu Laboratories  
of Europe, Ltd.  
(Europe)  
(Established 2001)



FORTUNE named Fujitsu as  
"one of the World's Most Admired Companies"  
for a third consecutive year.

A rich history  
**80 years**  
of Innovation



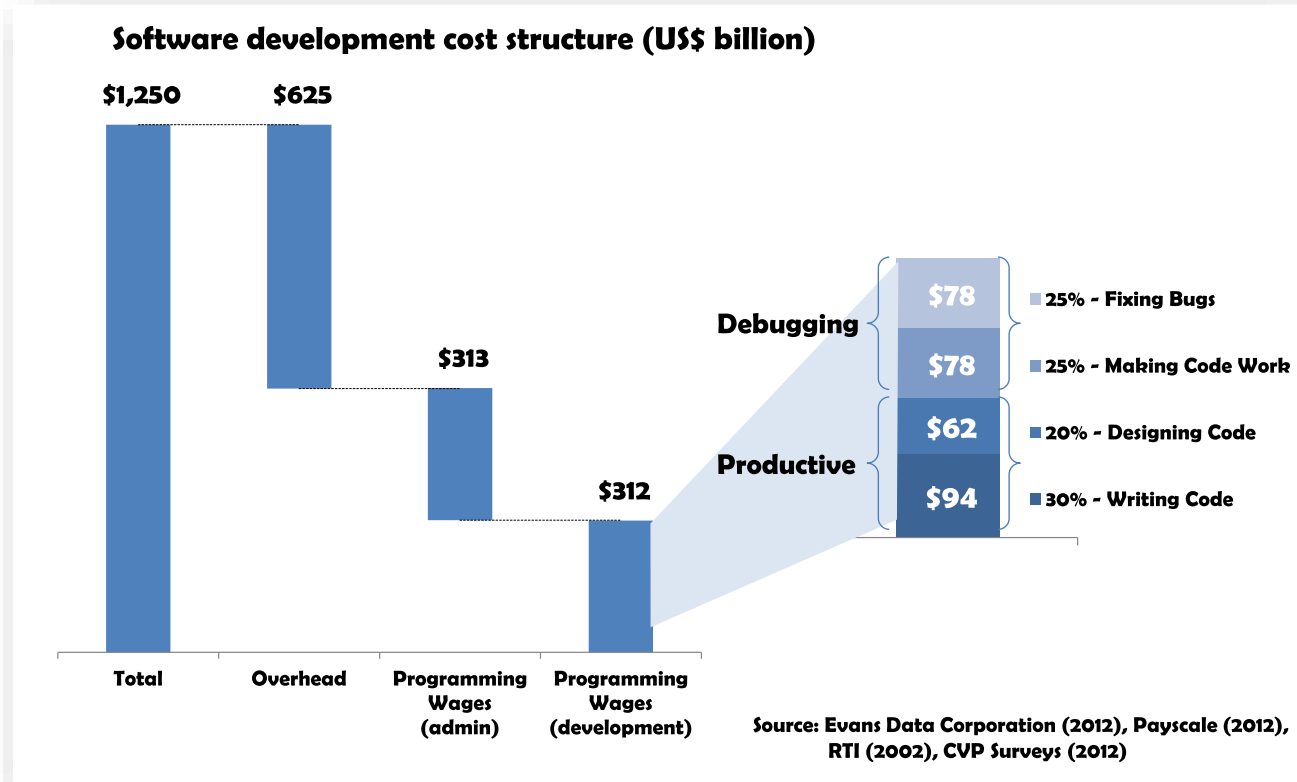
**156** thousand  
employees

On site service  
**180+**  
countries



# Motivation: Building Software is Expensive

Source: T. Britton et al., "Reversible Debugging Software," 2013.



The global cost of software development (as of 2013) was estimated at US \$1.25 Trillion !

- The annual global cost of debugging software is US \$312 billion
- Software developers spend *half their time* finding and fixing bugs

Tremendous business potential for software developer productivity enhancement.

# The Cost of (hardware & software) Bugs



$$\frac{4,195,835}{3,145,727} = 1.333820449136241002$$

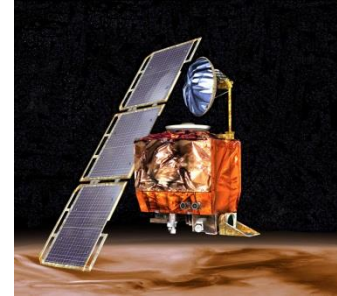
**1.333739068902037589**

## Intel Pentium FDIV Bug – 1994

**Total Cost: \$475 million**

## Mars Climate Orbiter- 1999

Destroyed due to software on the ground generating commands in pound-force (lbf), while the orbiter expected newtons (N).



**Mission Cost: \$327.6 million**

The global cost of debugging software (as of 2013) had risen to \$312 billion annually. The research found that, on average, software developers spend 50% of their programming time finding and fixing bugs.

- Cambridge University Research Study[1]

[1] <http://www.prweb.com/releases/2013/1/prweb10298185.htm>

## Heartbleed OpenSSL Vulnerability

**April 2014**



“The Heartbleed bug has likely cost businesses tens of millions of dollars in lost productivity” (Reuters)

# Software Testing $\Rightarrow$ Software Quality

- Testing is a dominant element in the processes to establish confidence in the correctness of software
- High-quality test suites are notoriously laborious to develop
  - **Product code of SQLite** version 3.7.17 consists of **81.3 KSLOC**, while its **test suite** is **91,421.1 KSLOC**, *i.e.*, **1,124x larger** than the product code itself
- Automatic test generation holds the promise of helping reduce the cost of software testing

# What kind of Software is relevant?

- Did an extensive survey within Fujitsu companies and subsidiaries

- This type of data is hard to come by not well documented

- Three major categories emerged

- Enterprise software mainly client-server type

- Server side code written in Java

- We have research and a tool on this based on Java Pathfinder

- Embedded software in networking switches, scanners, vending machines etc.

- Written in C/C++

- This will be a target of this talk

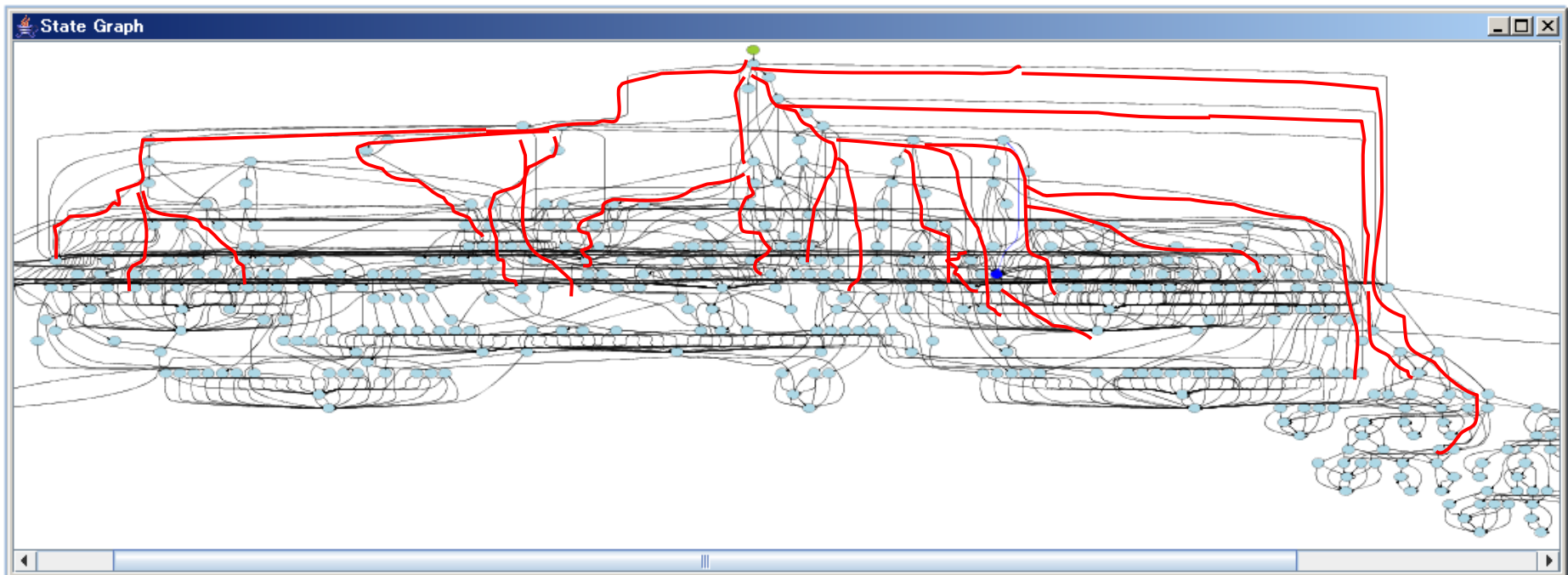
- Client/user side code in an UI/Browser

- Written in JavaScript, Python, Php etc.

- We have done some work on this too.

# Main problems in the QA process?

- Testing is time consuming and boring
  - Writing tests was a completely manual process
- Testing was often outsourced to third parties
  - Third parties do not have complete understanding of spec to test
- Test coverage is often poor





# Automated Test Generation

## 4. Test Flow Integration

```
class myClass
{
    myClass() {}
    int foo(int x)
    {
        if(x > 5)
            return 2;
        else
            return 0;
    };
};
```

**Function Under Test**

```
void TestDriver()
{
    myClass m;
    int test = 0;
    m.foo(test);
}
```

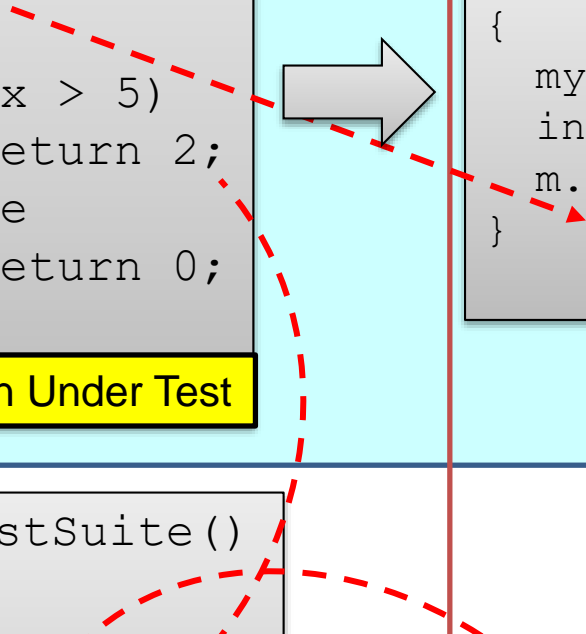
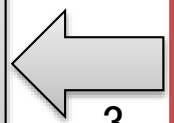
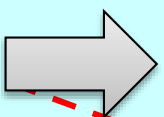
**1. Test Driver Generation**

```
void cppUnitTestSuite()
{
    myClass m;
    assert(m.foo(6) == 2);
    assert(m.foo(0) == 0);
}
```

**3. Test Oracle Generation**

```
Test vector 1: test = 0x0006
Test vector 2: test = 0x0000
```

**2. Test Input Generation**



## ■ Input generation

- Symbolic execution based test generation (KLEE based)
  - Implementing a String based SMT solver
- Unit level automatic test generation

## ■ Test-driver generation

- Iterative diagnosis-driven refinement

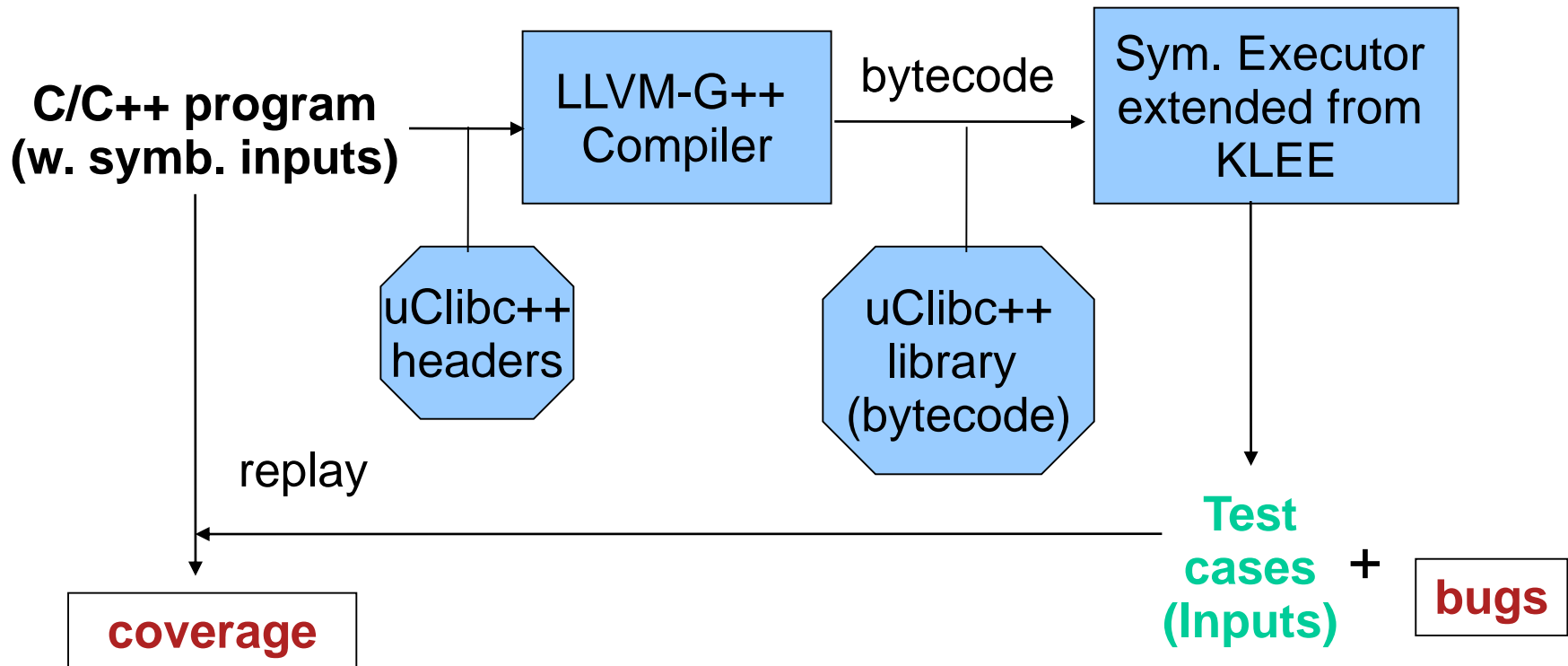
## ■ Oracle problem

- User defined, mutation, general errors or golden model

## ■ Test flow integration

- Incremental test generation from existing test-suites

# KLOVER: KLEE for C++ programs



- Extension of KLEE
- Add support for extra LLVM instructions, mainly intrinsic functions
  - >30 such instructions *e.g. `llvm.stacksave`, `llvm.memcpy`, ...*
  - C++ specific instructions, e.g. instructions (`llvm.atomic.*`) pertaining to the C++ memory model
- **Dynamically link** the bytecode of a C++ library (implementation) in runtime
  - we extend the publicly available library **uClibc++**
  - we optimize this library for symbolic execution

- Sometimes **slow**, and path **blow up**
  - It executes the low level bytecode for each API in the library
  - It relies on a constraint solver to calculate symbolic values (computations in AES are too complicated to the solver)
  - The implementation of an API may contain too many branches in the library implementation
    - Take the **string** library as example. Klee may generate thousands of paths for a simple string program, most of which are **useless** paths

# Our Optimizations

- To speed up KLOVER and scale it to large programs
  - Re-implement some core API in the uClibc++ library
    - Remove the branches within an API implementation so as to avoid useless paths
    - Intercept the calls to some APIs and replace the calls with customized handling
  - Develop our own solvers for constraint solving
    - The PASS string solver in FLA
    - Application-oriented solvers (e.g. for database applications)

# Eg: Original Implementation of string.compare (uClibc++)

```
_UCXXEXPORT int compare(const basic_string& str) const{  
    ...  
    int retval = strcmp(vector<Ch, A>::data, str.vector<Ch, A>::data, len);  
    ...  
}
```

```
int strncmp(register const char *s1,  
            register const char *s2, size_t n) {  
    while (n && (*s1 == *s2)) {  
        if (!*s1++)  
            return 0;  
        ++s2;  
        --n;  
    }  
    return (n == 0) ? 0 : (*s1 < *s2) ? -1 : 1;  
}
```

KLOVER will  
create  $> n$   
branches within  
this API call

Two calls to this API  
will lead to  $> n^2$   
paths

# Example: Re-Implementation of string.compare

```
_UCXXEXPORT int compare(const basic_string& str) const {  
...          // some initialization work  
int v = 0;    // 1, 0 and -1 stand for gt, eq and lt respectively  
  
for (size_type i = 0; i < rlen; i++) {  
    v += (!v) * ((operator[](i) > str[i]) - (operator[](i) < str[i]));  
}  
  
v += (!v) * ((vector<Ch, A>::elements > str.elements) -  
    (vector<Ch, A>::elements < str.elements));  
return v;  
}
```

KLOVER now explores only **one** path within this API call; no useless branches will be created (note that no “if” or “while” is involved).



# More Issues

- Executing the API code is slow even using the improved implementation
- It cannot handle symbolic lengths

```
for (size_type i = 0; i < rlen; i++) {  
    ...  
}
```

KLOVER will fail  
when rlen is  
symbolic



- We prefer a more **capable** and **efficient** solver
  - We extend the KLEE **IR** to support string operations
  - We intercept all string API calls and build the IR
  - We check the satisfiability of a string IR expression with **customized solver** which support **symbolic lengths**

# Handling Symbolic Strings

```
foo(String a, String b) {  
  String c;  
  String d;  
  c = a.concat(b);  
  if !(c.equals("qrs")) {  
    d = c.concat("t");  
    return d;  
  }  
  return c;  
}
```

We need to create a solver that solves for such string constraints

$\Phi$  is symbolic expression  
 $x, y, z, w$  are symbolic strings

$a = x, b = y, c = z, d = w, \Phi = \{ \}$

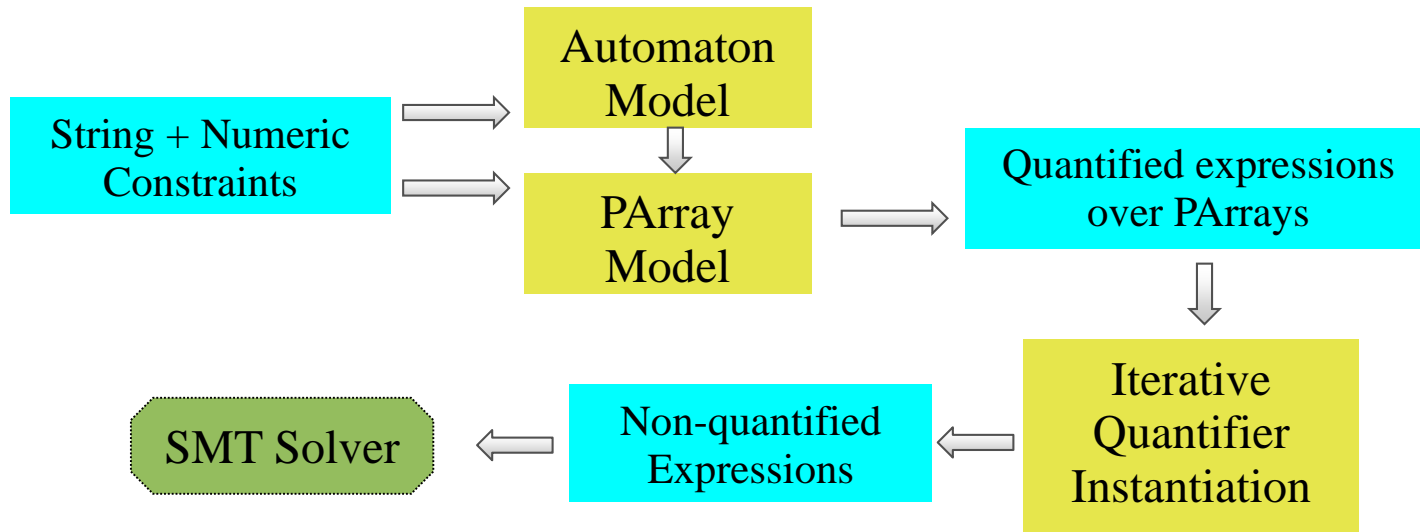
$\Phi = \{ z = \text{concat}(x, y) \}$

$\Phi = \{ (z = \text{concat}(x, y) \ \& \ (z \neq \text{"qrs"})) \}$

$\Phi = \{ (z = \text{concat}(x, y)) \ \& \ (z = \text{"qrs"}) \}$

$\Phi = \{ (w = \text{concat}(\text{concat}(x, y), \text{"t"})) \ \& \ (z \neq \text{"qrs"}) \}$

- Model strings as parameterized arrays (PArrays)
  - $s1 : [0][1] \dots [len-1]$
  - both the index and the length are **parameterized** (i.e. **symbolic**)
  - string constraints are **quantified expressions** with symbolic indices
    - e.g.  $\forall i \in [k, k+10] : s1[i] \neq c$ , for symbolic  $k$  and  $c$
- Use **quantifier elimination** to find solutions or prove unsat
  - use symbolic length based heuristics to remove quantifiers
  - always terminate for bounded lengths





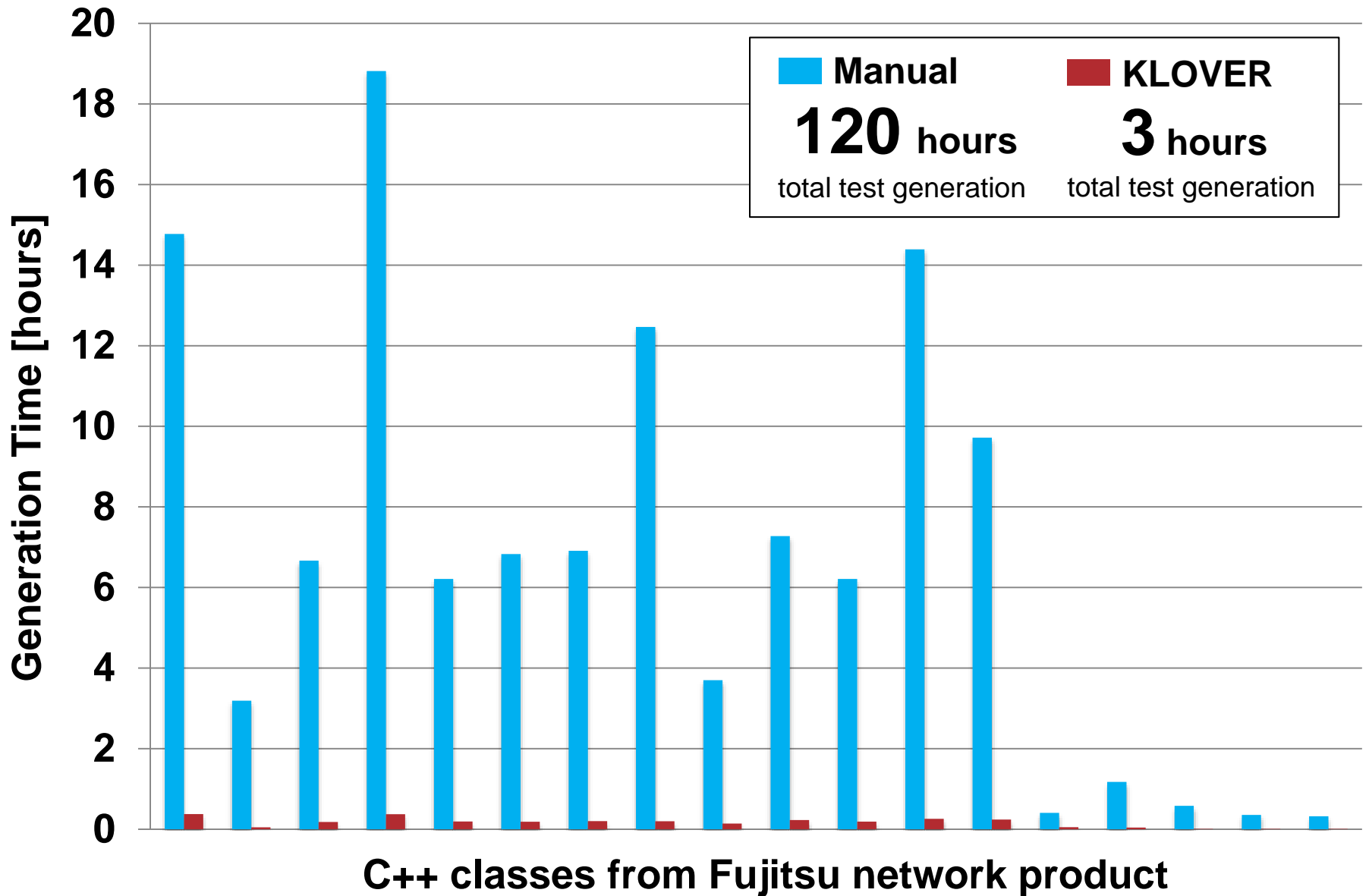
4 million lines of C++ code

- Manual testing: 4hrs per class
- Unit test quality was poor

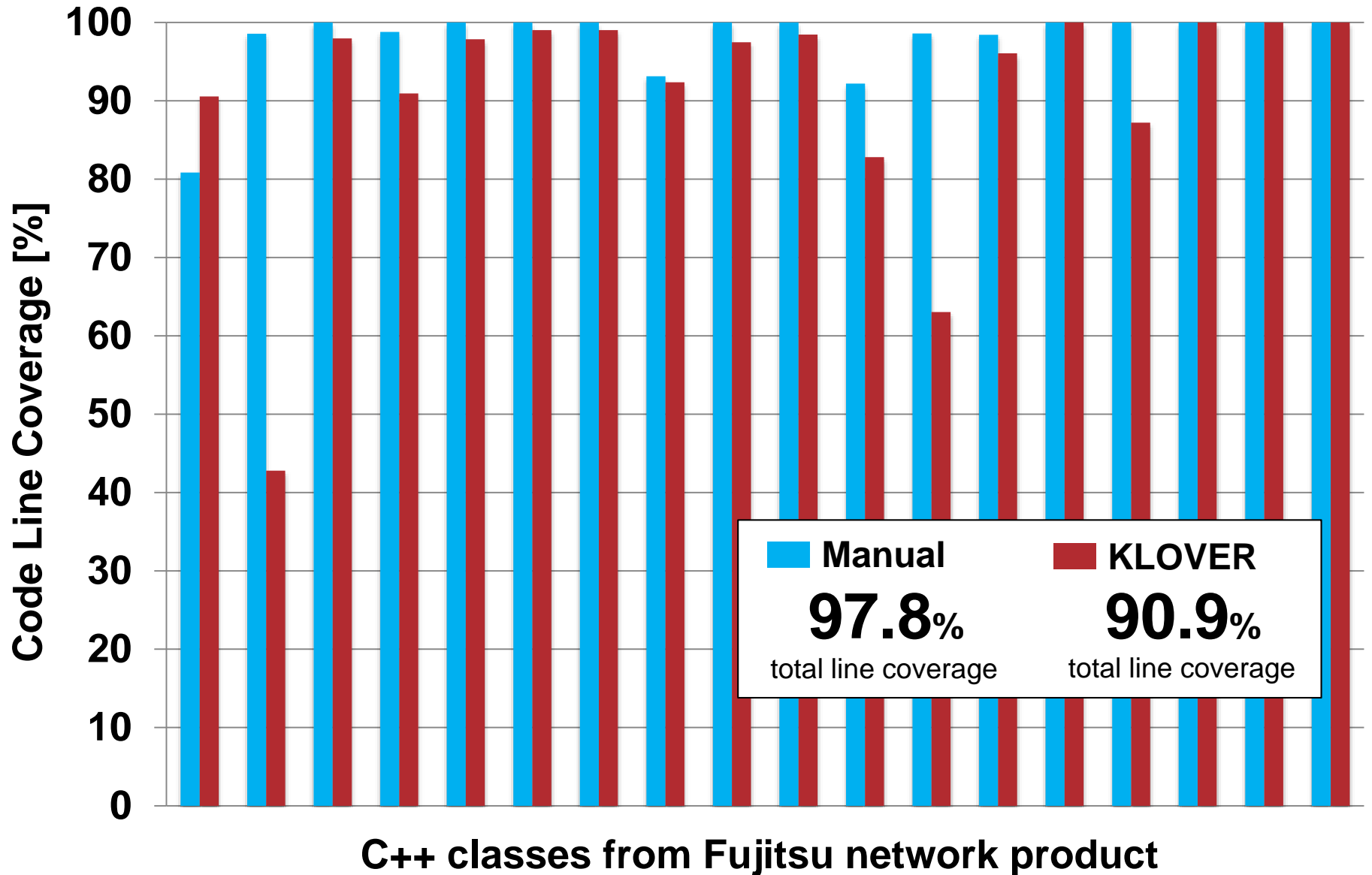
## FUJITSU LAYER 2 PHOTONIC SWITCH

- ❑ Applied KLOVER to 5 of 50 modules in switch code
  - ❑ Total about 500K lines of code
- ❑ KLOVER achieved 80% line coverage and found several bugs in code
- ❑ Technology was subsequently transferred to Fujitsu Network Computing, Texas

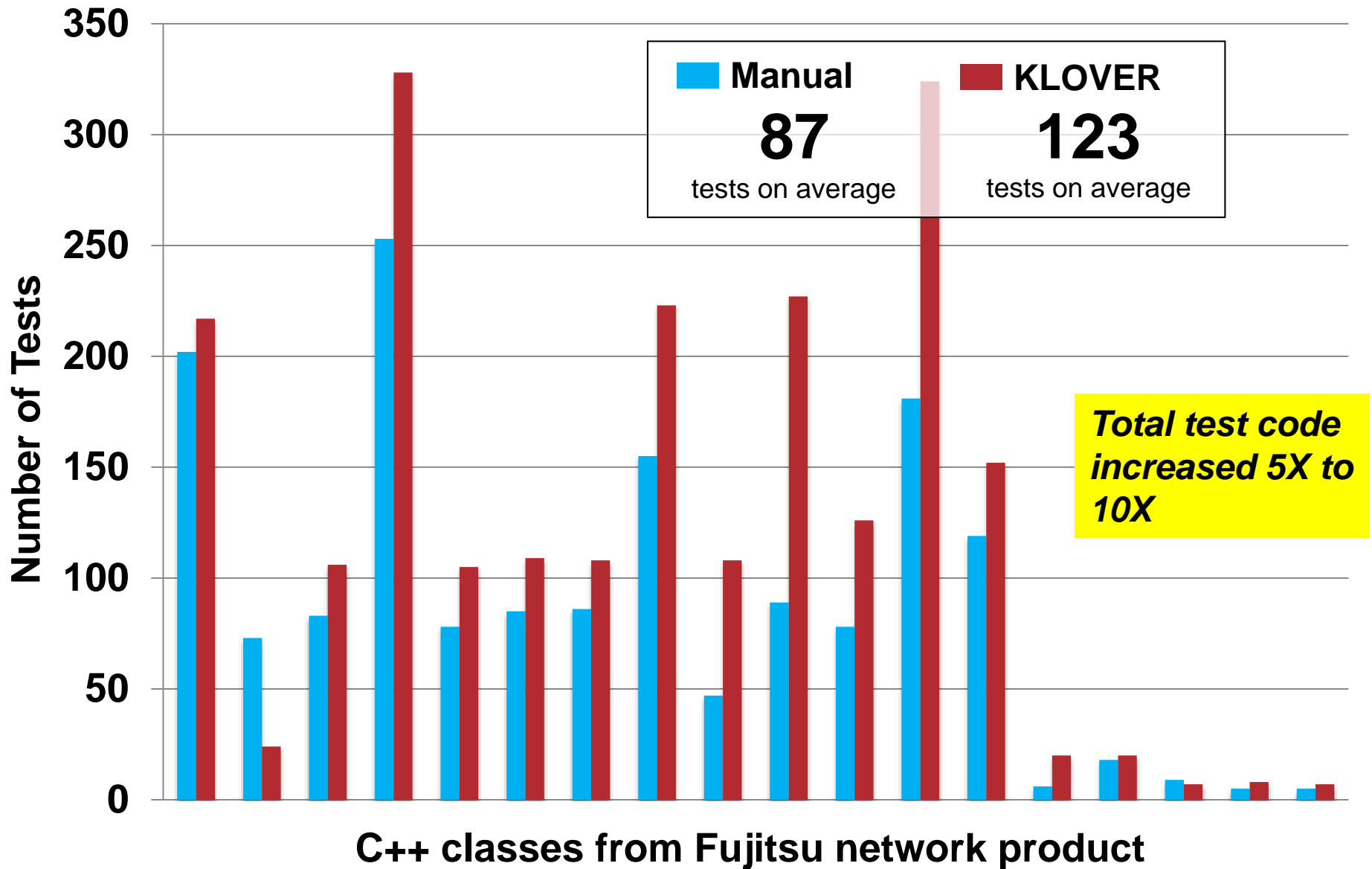
# Benchmark Results: Test Generation Time



# Benchmark Result: Code Coverage



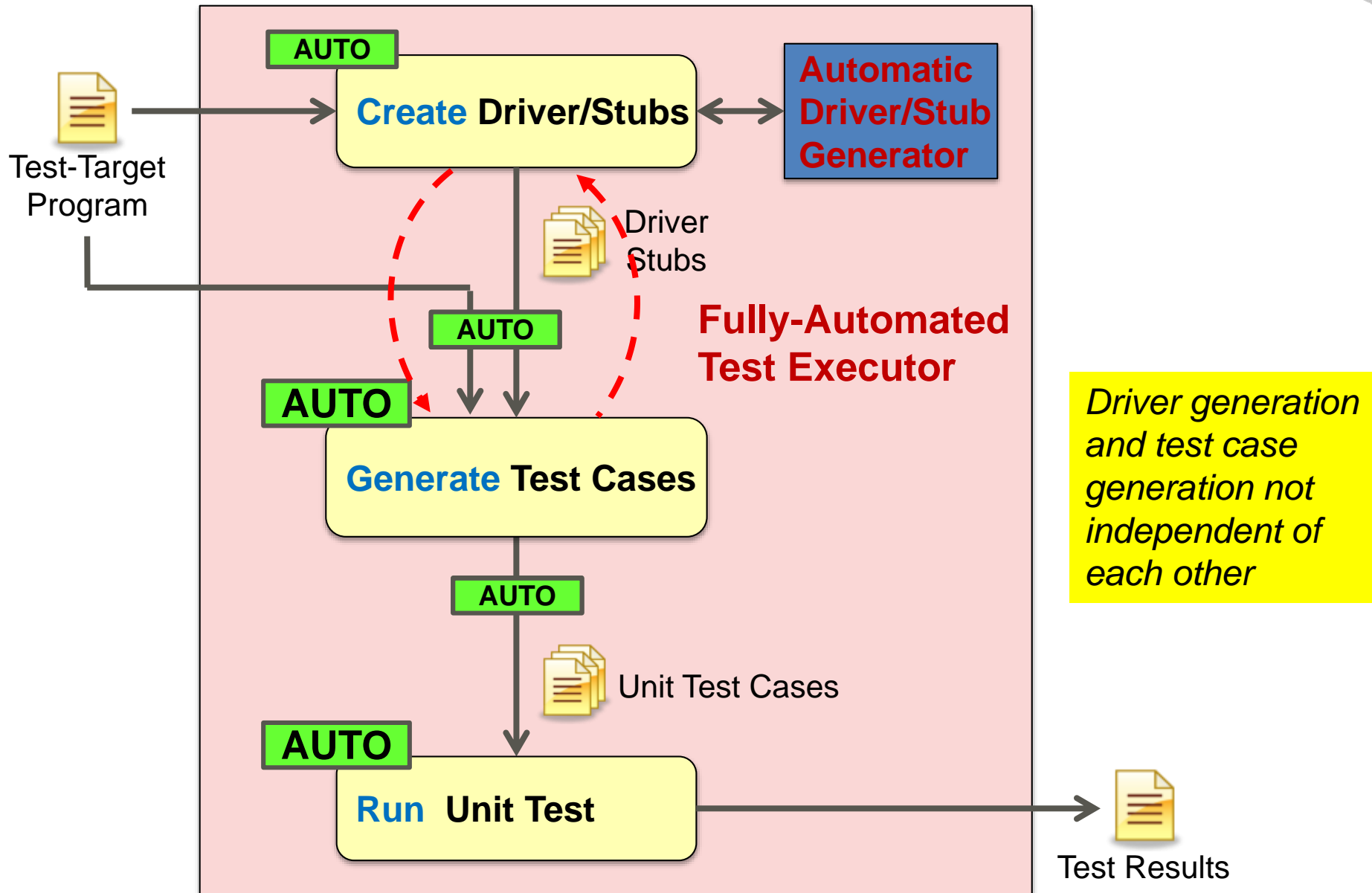
# Benchmark Result: Number of Tests



- Huge number of new test cases generated
- Difficult to understand intent of new tests
- Exceptions caused by unconstrained/invalid inputs
- Difficult to maintain such huge test suite
- The importance of writing compact test cases is well recognized
  - Quotes from the GNU bug reporting instructions (LLVM, Mozilla & Webkit have similar guidelines):
    - *“smaller test cases make debugging easier”*
    - *“GCC developers prefer bug reports with small, portable test cases”*
    - *“minimized test cases can be added to the GCC test suites”*



# Test Case quality depends on Test Driver quality



# Motivational Example: Naïve Approach

```
void naive_driver() {  
    void* x = malloc(1024);  
    register_symbolic(x);  
    int exp;  
    register_symbolic(exp);  
    assert(packet_receive(x) == exp);  
}
```

## Naïve Driver

```
int packet_receive(void* x) {  
    packet* y = (packet*) x;  
    memcpy(y->payload, y->buf, y->size);  
    data* z = (data*) y->payload;  
    if (z->flag == FLAG_OK) {  
        return 0;  
    }  
    return 1;  
}
```

Not obvious what the function expects. Pass a symbolic memory object of some size

`packet_receive()` implicitly assumes

- `x` points to `packet` object
- `y->size` is the size of `data` object
- `y->payload` points to `data` object

Symbolic execution would not be able to cover this branch

*Intelligent driver generation is the key!*

# Example: Diagnosis-Driven Driver Refinement

```
void driver_1() {  
    void* x = NULL;  
    int exp;  
    register_symbolic(exp);  
    assert(packet_receive(x) == exp);  
}
```

**Minimal Driver**

```
void driver_2() {  
    data* x_buf = new data();  
    register_symbolic(x_buf->flag);  
    packet* x = new packet();  
    x->size = sizeof(data);  
    x->buf = x_buf;  
    x->payload = new data();  
    int exp;  
    register_symbolic(exp);  
    assert(packet_receive(x) == exp);  
}
```

**Enhanced Driver**

Symbolic Execution

```
int packet_receive(void* x) {  
    packet* y = (packet*) x;  
    memcpy(y->payload, y->buf, y->size);  
    data* z = (data*) y->payload;  
    if (z->flag == FLAG_OK) {  
        return 0;  
    }  
    return 1;  
}
```

## Diagnosis Information

- (1) **x** is accessed as **packet** type
- (2) **sizeof(y->payload) ≥ y->size**
- (3) **y->payload** is accessed as **data** type
- (4) **z->flag** field is used in path condition

Driver Enhancement

# Example: Iterative Test Generation

- Re-uses existing test-cases from the previous iteration, cloning and modifying them minimally at fine-grained level to create new test-cases

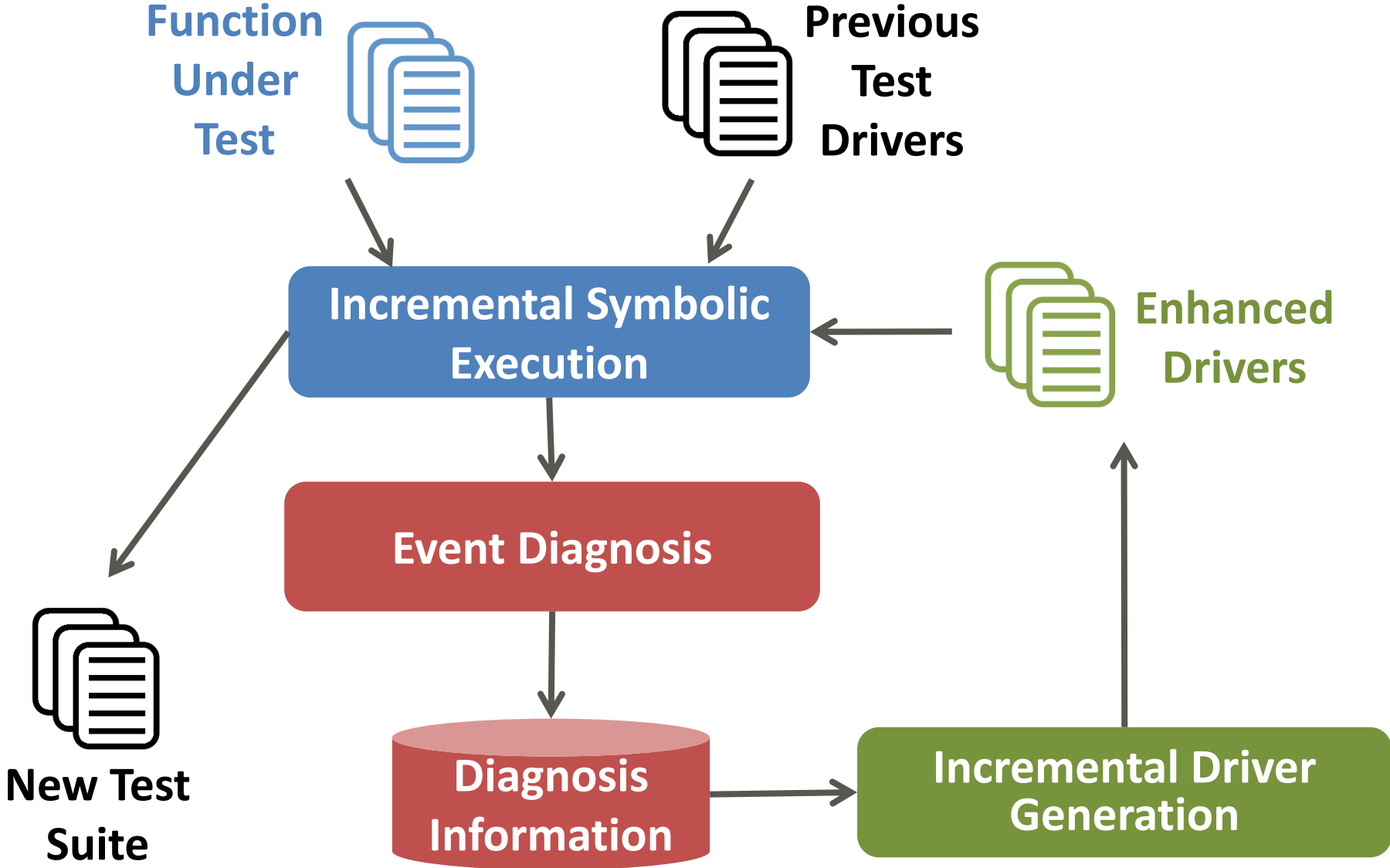
```
void test_1() {  
    data* x_buf = new data();  
    x_buf->flag = FLAG_OK;  
    packet* x = new packet();  
    x->size = sizeof(data);  
    x->buf = x_buf;  
    x->payload = new data();  
    assert(packet_receive(x) == 0);  
}
```

1<sup>st</sup> iteration generates **test\_1()**

2<sup>nd</sup> iteration modifies two lines  
with an annotation "**// !test\_1**".

```
void test_2() {  
    data* x_buf = new data();  
    x_buf->flag = FLAG_INVALID; // !test_1  
    packet* x = new packet();  
    x->size = sizeof(data);  
    x->buf = x_buf;  
    x->payload = new data();  
    assert(packet_receive(x) == 1); // !test_1  
}
```

# Overview: Fine-Grained Incremental Generation



# Incremental Symbolic Execution

	Basic Symbolic Execution	C++ Support	Incremental Execution	Event Diagnosis
<b>KLEE</b>	✓			
<b>KLOVER</b>	✓	✓		
<b>FSX</b>	✓	✓	✓ <i>(Similar to DiSE &amp; Memoise)</i>	✓

*Note: FSX tool is built from scratch.  
It does not share code with any of these tools.*

*Outputs a diagnosis for a given indicative event*

```
int packet_receive(void* x) {  
    packet* y = (packet*) x;  
    memcpy(y->payload, y->buf, y->size);  
    data* z = (data*) y->payload;  
    if (z->flag == FLAG_OK) {  
        return 0;  
    }  
    return 1;  
}
```

## Indicative Events

- (1) Null pointer access at `y->payload`
- (2) Out-of-bound access at `memcpy()`
- (3) Branch-not-taken at `if (z->flag == FLAG_OK)`

## Type Information

`y` is accessed as **packet** type  
`z` is accessed as **data** type

## Range Information

`sizeof(y->payload) ≥ y->size`

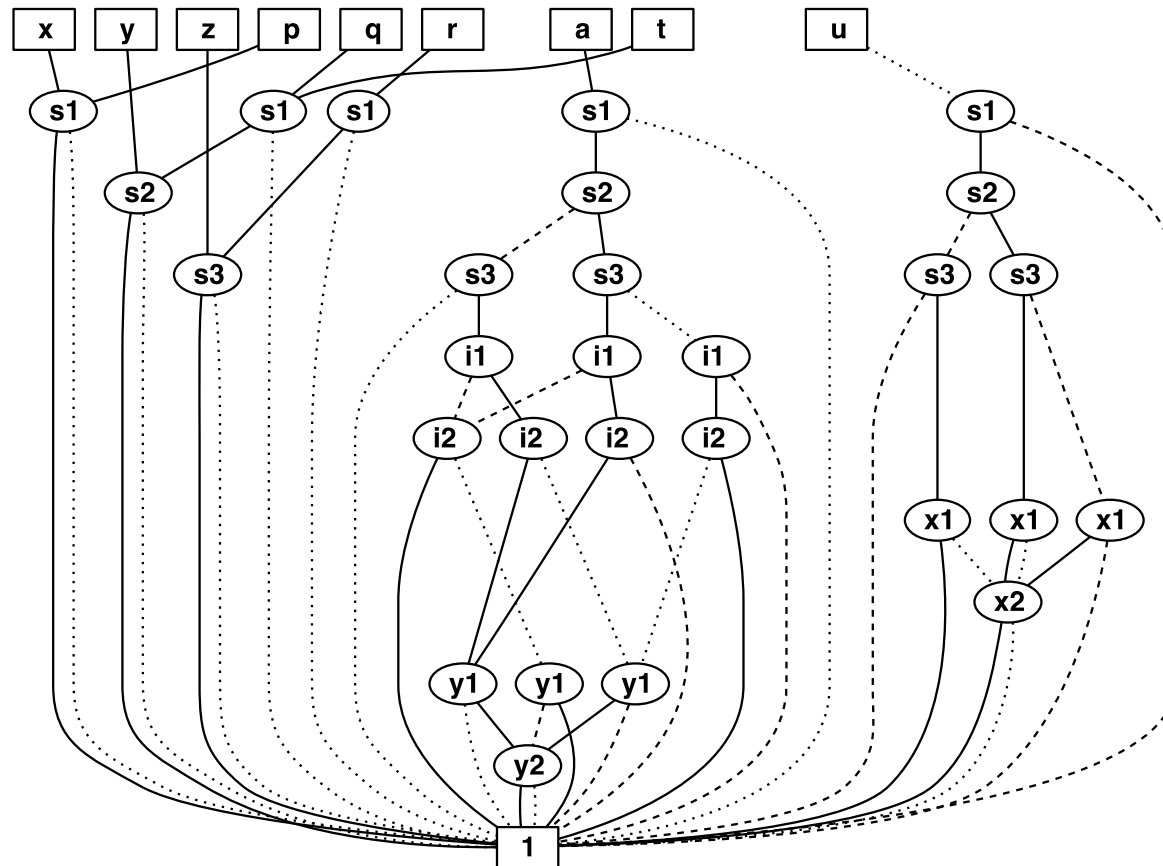
## Relevant Input Set

`y → x, y->size → x->size`  
`y->buf[] → x->payload[]`  
`z->flag → x->payload->flag`

Diagnosis Information

# Precise Computation of Relevant Input Sets

*Encoded as Boolean characteristic functions, and compactly represented through Reduced-Ordered Binary Decision Diagrams (ROBDDs)*





# Incremental Driver Generation

*Enhances previous drivers to eliminate events by modifying values of inputs relevant to such events*

## Null pointer access

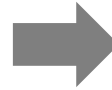
```
void* x = NULL;  
.  
.  
.  
memcpy(y->payload, y->buf, y->size);
```



```
void* x = (void*) new packet();  
.  
.  
.  
memcpy(y->payload, y->buf, y->size);
```

## Out-of-bound access

```
x->size = 0;  
.  
.  
.  
memcpy(y->payload, y->buf, y->size);
```



```
x->size = sizeof(data);  
.  
.  
.  
memcpy(y->payload, y->buf, y->size);
```

## Branch-not-taken

```
x->buf->flag = FLAG_DEFAULT;  
.  
.  
.  
if (z->flag == FLAG_OK) {
```



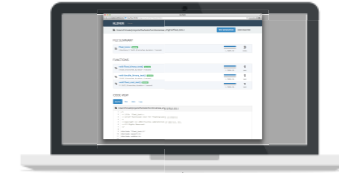
```
register_symbolic(x->buf->flag);  
.  
.  
.  
if (z->flag == FLAG_OK) {
```

# Tool Implementation

Command Line UI



Web UI



CodeCasa SCCS & CI Platform (Ruby)



FSX Server (Ruby)

FSX Test Generation Driver (Ruby)

FSX Core Engine (45k LOC in C++)

Incremental Driver Generator

Incremental Symbolic Executor

Relevant Input Analyzer

Mutation Analyzer

Test Stub Generator

Clang C/C++ Frontend

LLVM IR

CUDD Package

SQLite3 Database

Customized STP Constraint Solver

## ■ Baseline tool: FSX-Baseline

- Generates naïve driver where all assignable variables including function arguments, member variables and global variables are assigned symbolic values and any pointers are set to new objects
- Symbolic executor implements same search strategy as KLEE

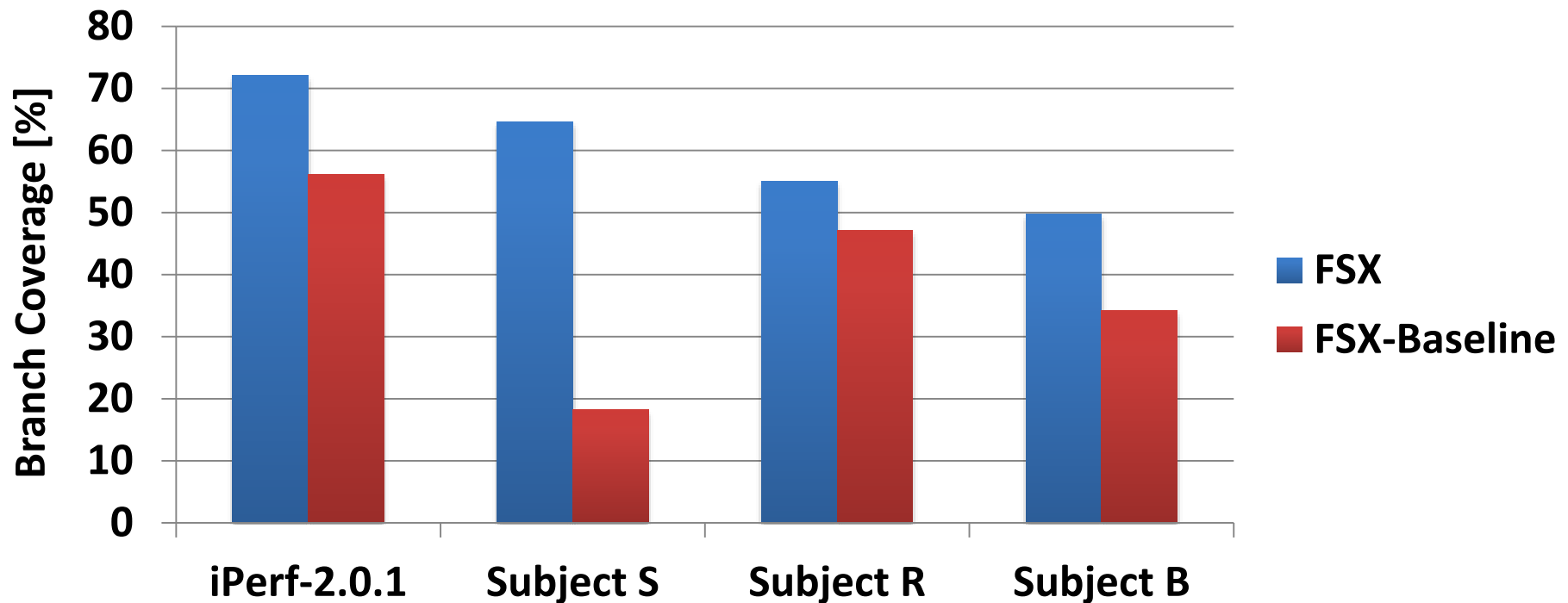
## ■ Benchmark subjects

- iPerf: network bandwidth measurement tool (5k LOC in C++)
  - Representative set of challenging issues for automatic test generation
  - 5 versions from version 2.0.1 to 2.0.5 used as a software evolution example
- Three embedded software subjects from Fujitsu commercial network products
  - Subject S (39k LOC in C), Subject R (12k LOC in C), Subject B (15k LOC in C)

# RQ1: Full Test Suite Generation

*Can FSX generate high-quality unit tests for large system software?*

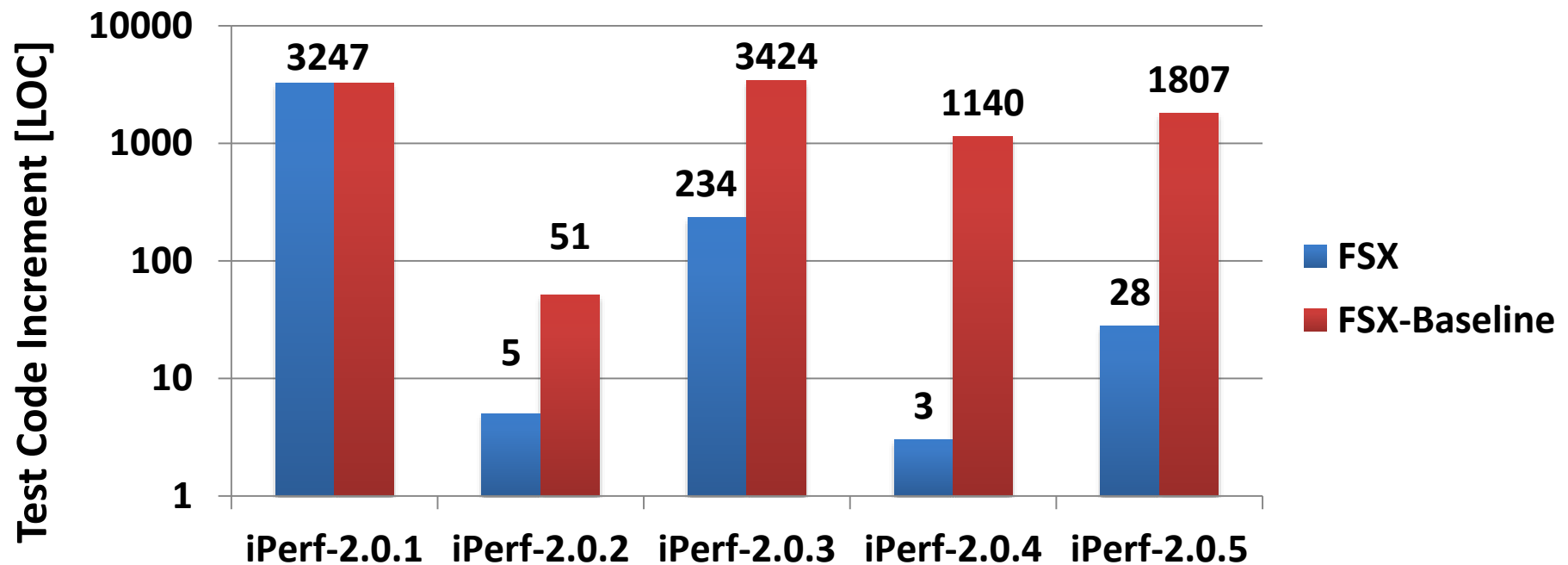
- FSX is able to generate higher quality tests than FSX-Baseline for all benchmarks



# RQ2: Test Suite Augmentation

*Can FSX perform test-suite augmentation, minimizing maintenance cost of new test code?*

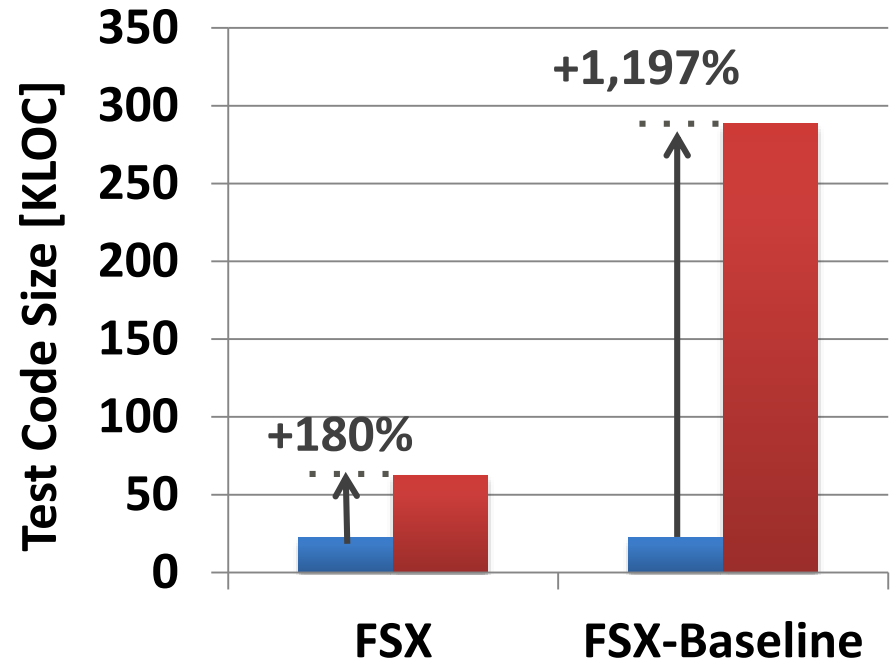
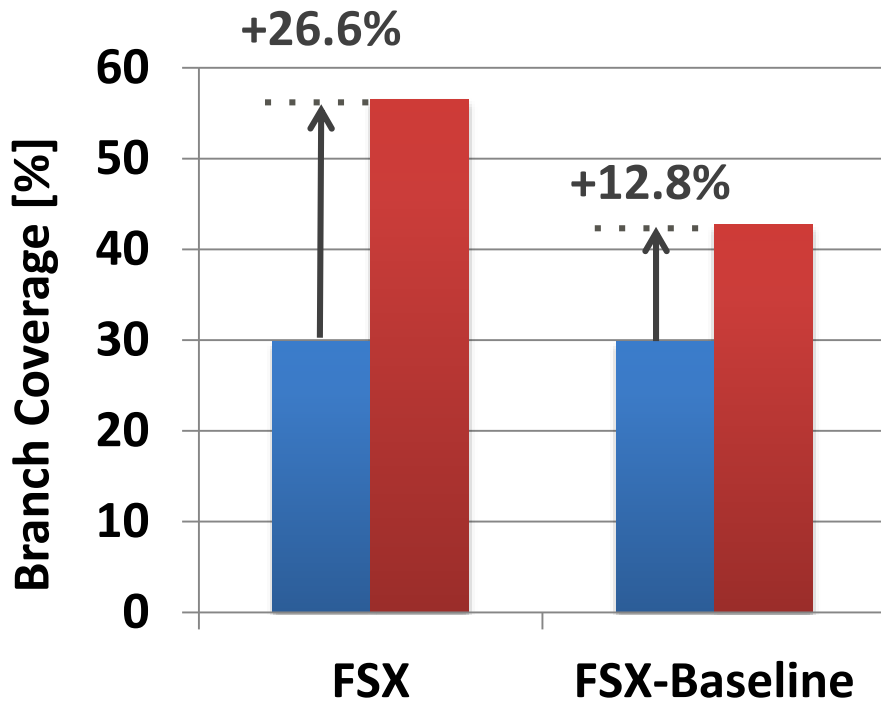
- FSX re-used existing test-cases from the previous version, cloning and modifying them minimally while FSX-Baseline generated new test-cases for uncovered functionality



# RQ3: Test Suite Enhancement

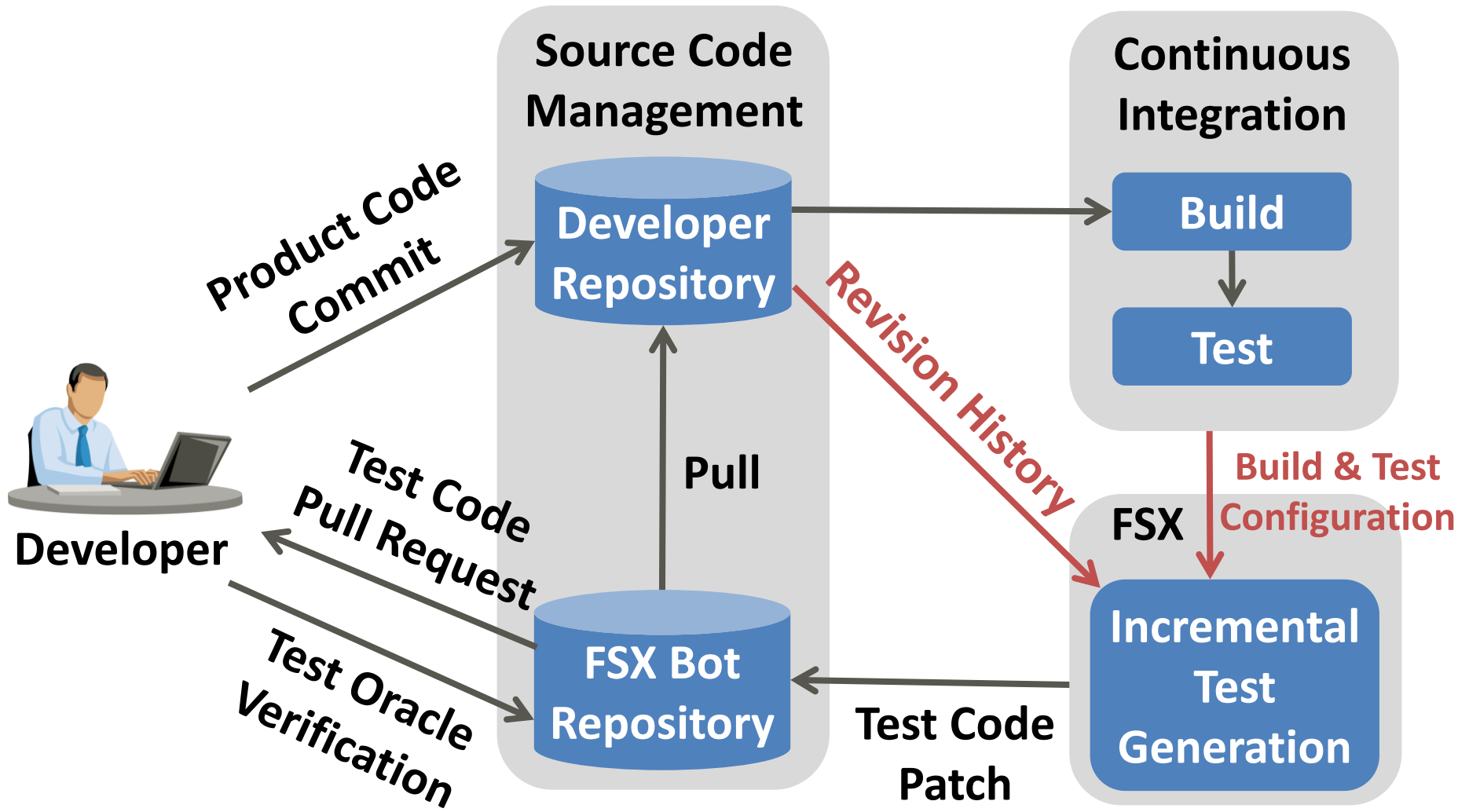
*Can FSX enhance an existing test-suite, minimizing maintenance cost of the new test code?*

- FSX's enhancement boosts the coverage much more than FSX-Baseline's while adding much less lines



- No easy answer
- Regression test suites from previous model
- User defined assertions in code
- Normal error conditions like exceptions, crashes etc.
- Automatic translation from a specification language like UML
- Using UI feature based invariants

# Use-case 1: SCM/CI System Integration





# Use-case 2: Application Modernization

Product Code  Test Code (Optional) 

Test Suite Generation

  
Generated Test Code

*Test oracles assert the current output value on legacy platform*

Do All Tests Pass?

NO

Bug Found!

YES

Complete



Legacy Platform



Modernization



Modern Platform

Product

## Buggy Product Code Example

```
if (memcmp(x, y, n) == -1) {  
    . . .  
}  
. . .
```

RHEL 5: `memcmp()`  
returns either -1, 0, or 1

RHEL 7: `memcmp()`  
returns the value  
difference between  
two memory blocks

Legacy Vending System

Modern Vending System

- Fujitsu research concentrated on making automated software testing useable in an industrial development environment
- Created techniques for test driver generation, test input generation, and seamless integration into the test and development cycle at unit testing level
- Tools are being used internally by software teams to enhance test and debug productivity
- Future Challenges:
  - scaling to larger modules for system or integration test
  - better ways of generating test oracles
  - better human – tool coordination
  - some effort in automatic debugging and repair

# Acknowledgements




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Hiroaki Yoshida



Guodong Li



**FUJITSU**

shaping tomorrow with you