

Fine-grain Memory Object Representation in Symbolic Execution

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SOFTWARE RELIABILITY
GROUP

2nd International



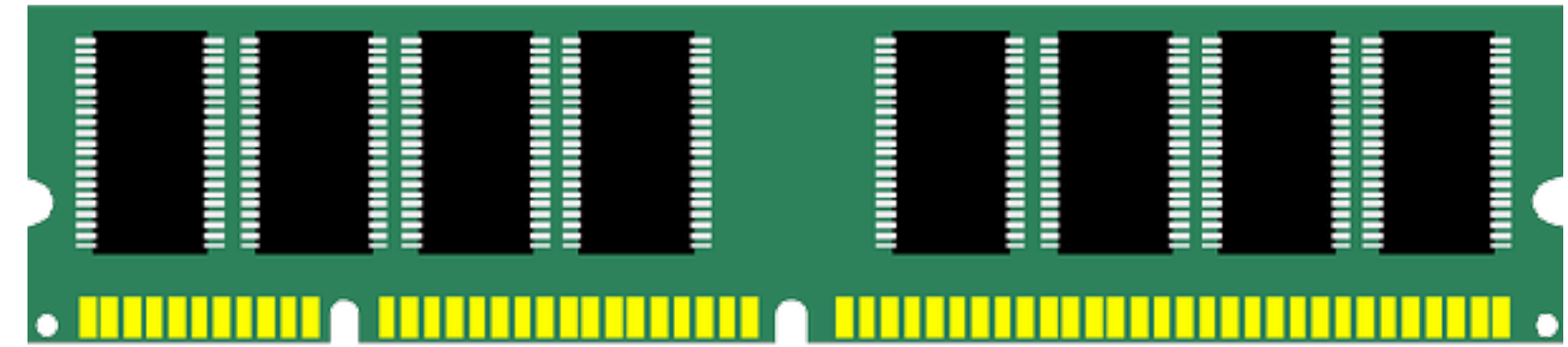
Workshop

Programs

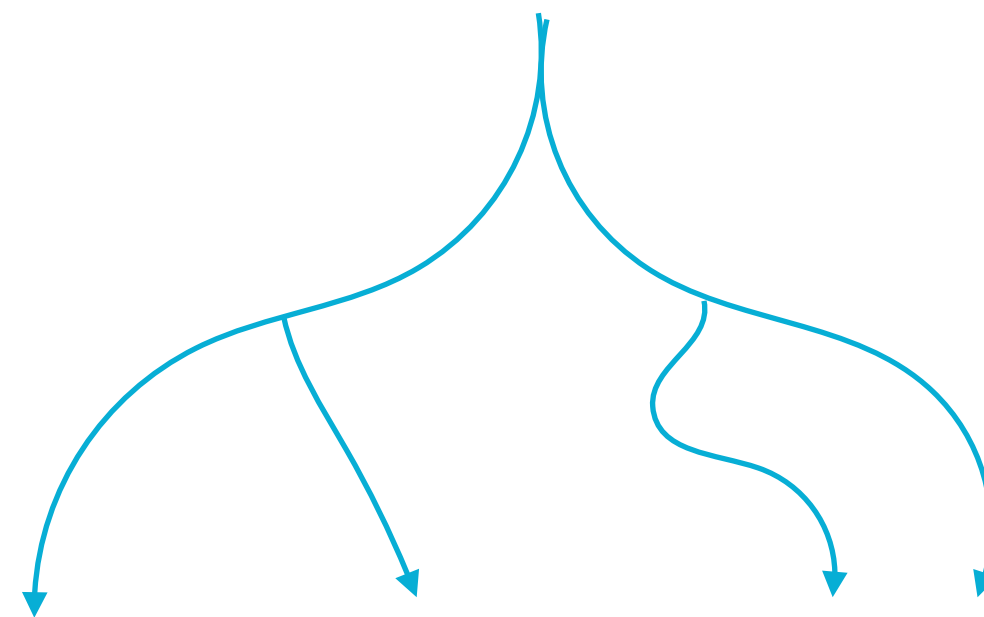
```
char * a = malloc(1024);  
int32 i = 10;
```

```
a[i]++;  
if (i != 12345)  
{  
    a[i-2] = a[i] * 2;  
} else {  
    a[i+2] = a[i] - 2;  
}
```

Memory Representation



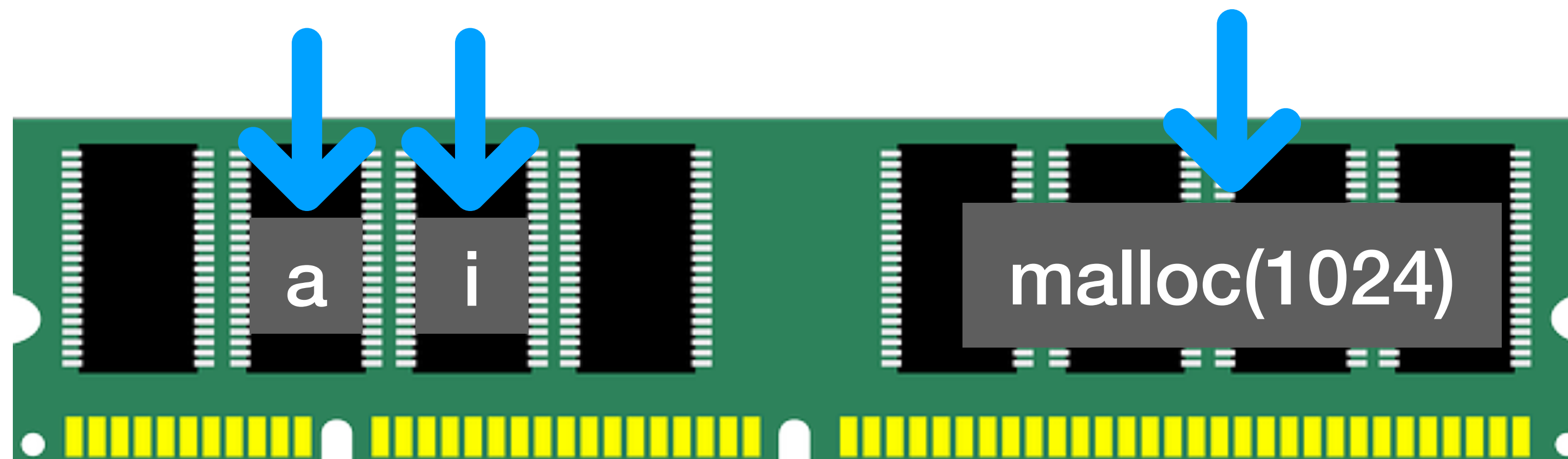
Symbolic Execution





```
char * a = malloc(1024);  
int32 i = 10;
```

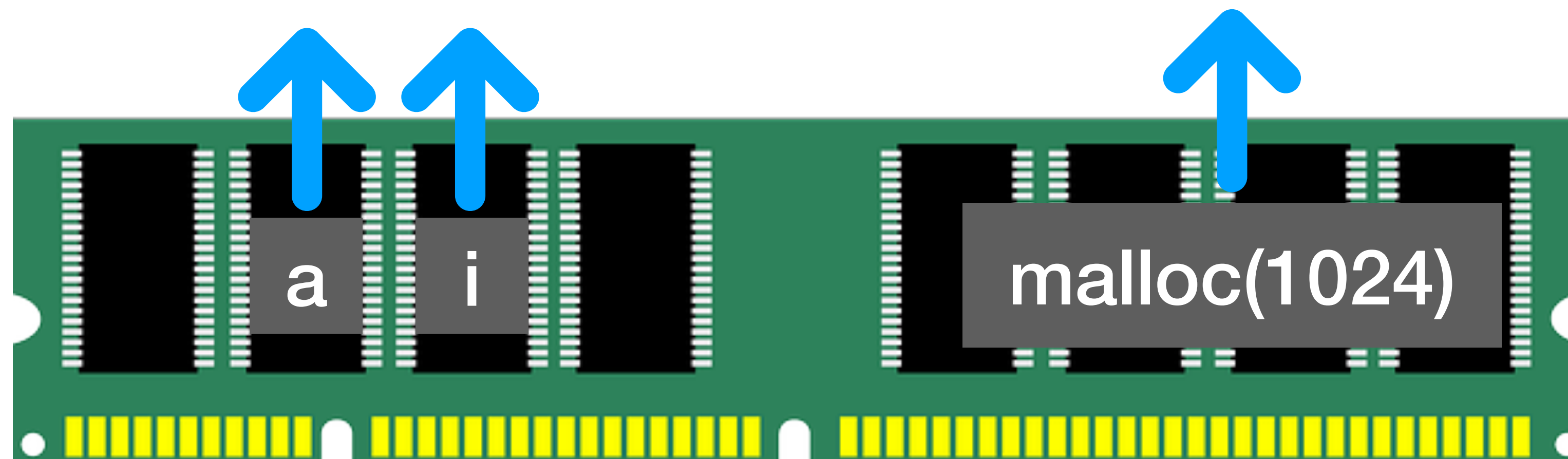
```
a[i]++;  
if (i != 12345) {  
    a[i-2] = a[i] * 2;  
} else {  
    a[i+2] = a[i] - 2;  
}
```





```
char * a = malloc(1024);  
int32 i = 10;
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a[i]++;  
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} else {  
    a[i+2] = a[i] - 2;  
}
```



(Dynamic) Symbolic Execution

```
char * a = malloc(1024);  
int32 i = symbolic;
```

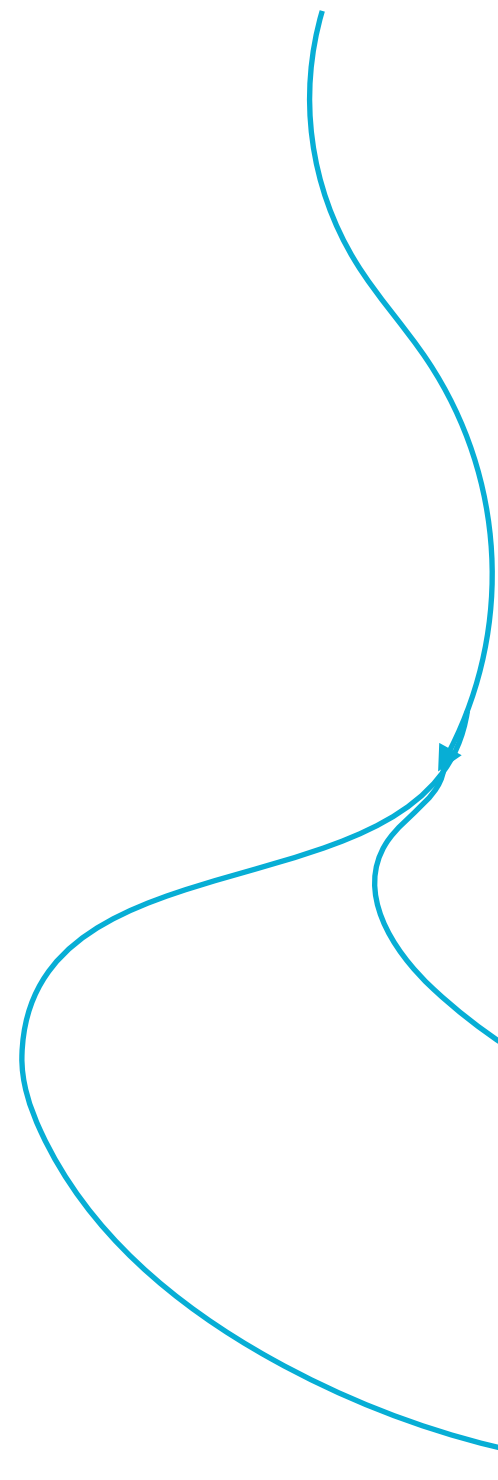
```
a[i]++;  
if (i != 12345)  
{
```

```
    a[i-2] = a[i] * 2;
```

```
    } else {
```

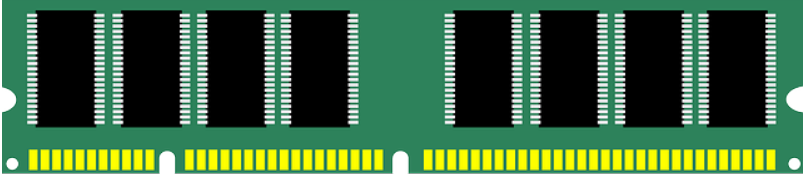
```
        a[i+2] = a[i] - 2;
```

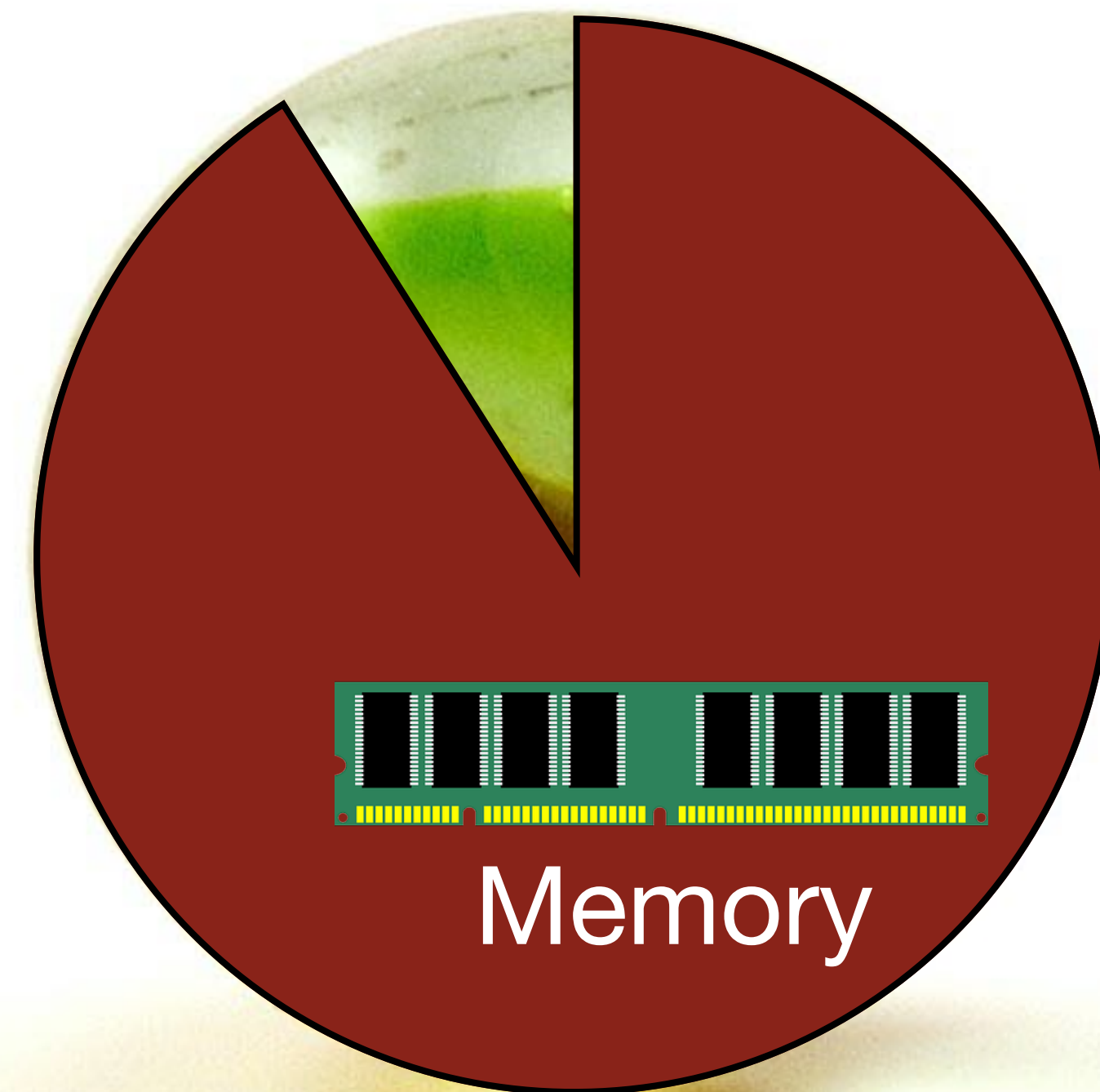
```
    }
```



Application State

A Simplified View

- Path Constraints
- Registers (i.e., program counter)
- Allocated Memory 
 - Stack
 - Heap

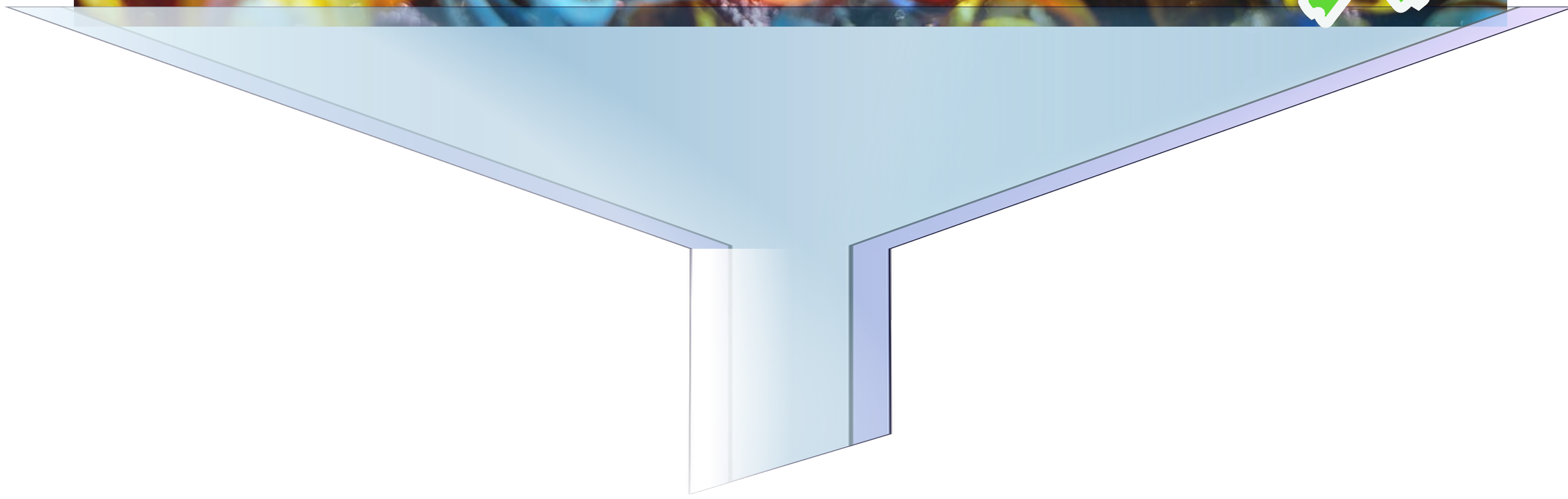


The many states ...

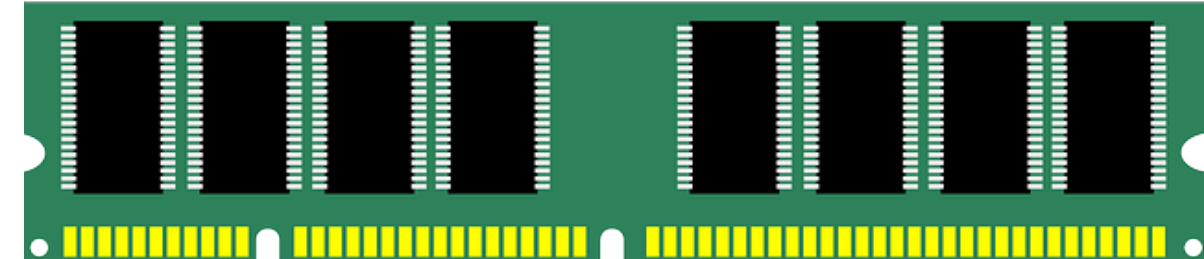


The many states ...





KEE



Handling Allocations

State of the art

Copy on Write (Cow)



malloc(1024)



```
char * a = malloc(1024);
```

```
...
```

```
if (i != 12345)
```

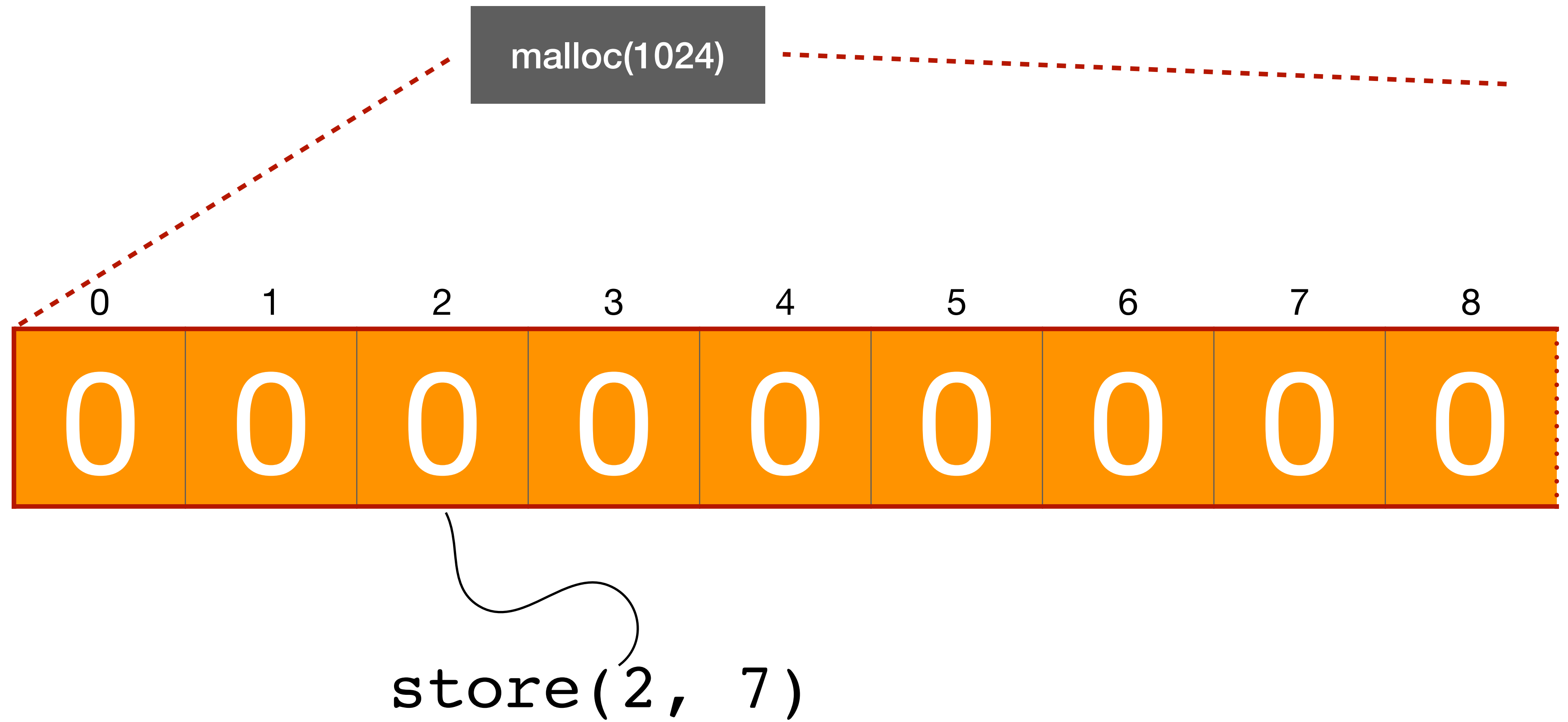
```
{
```

```
...
```

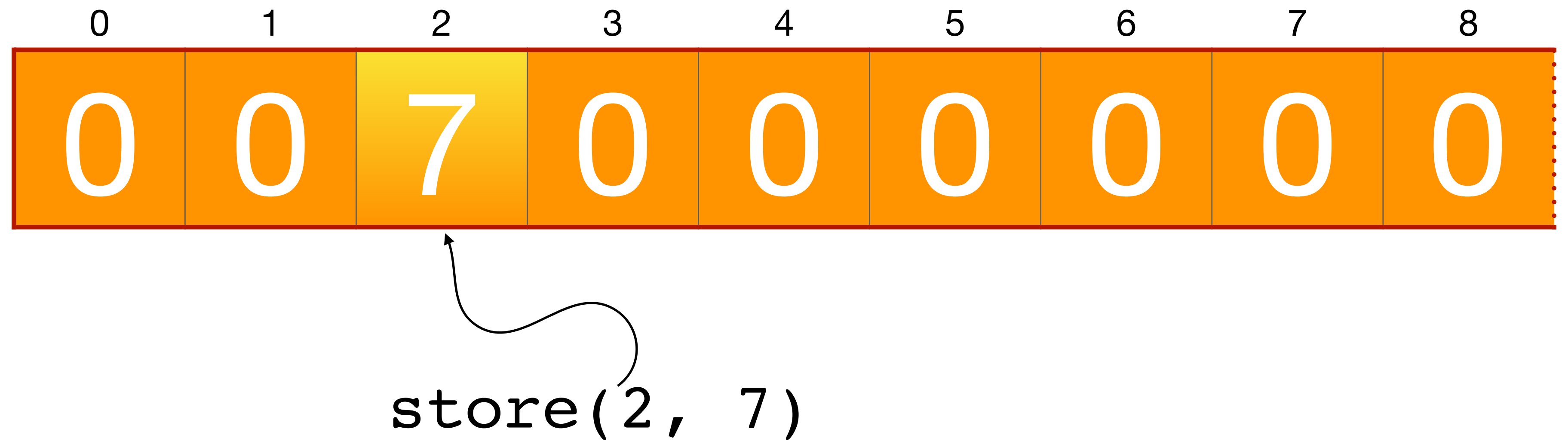
```
} else {
```

```
...
```

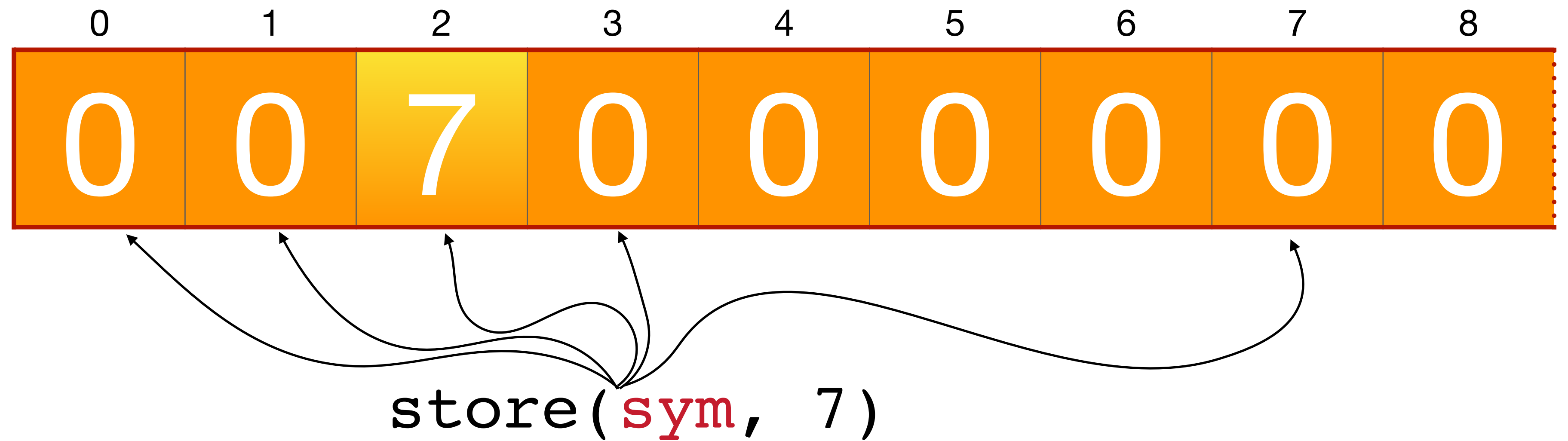
```
}
```



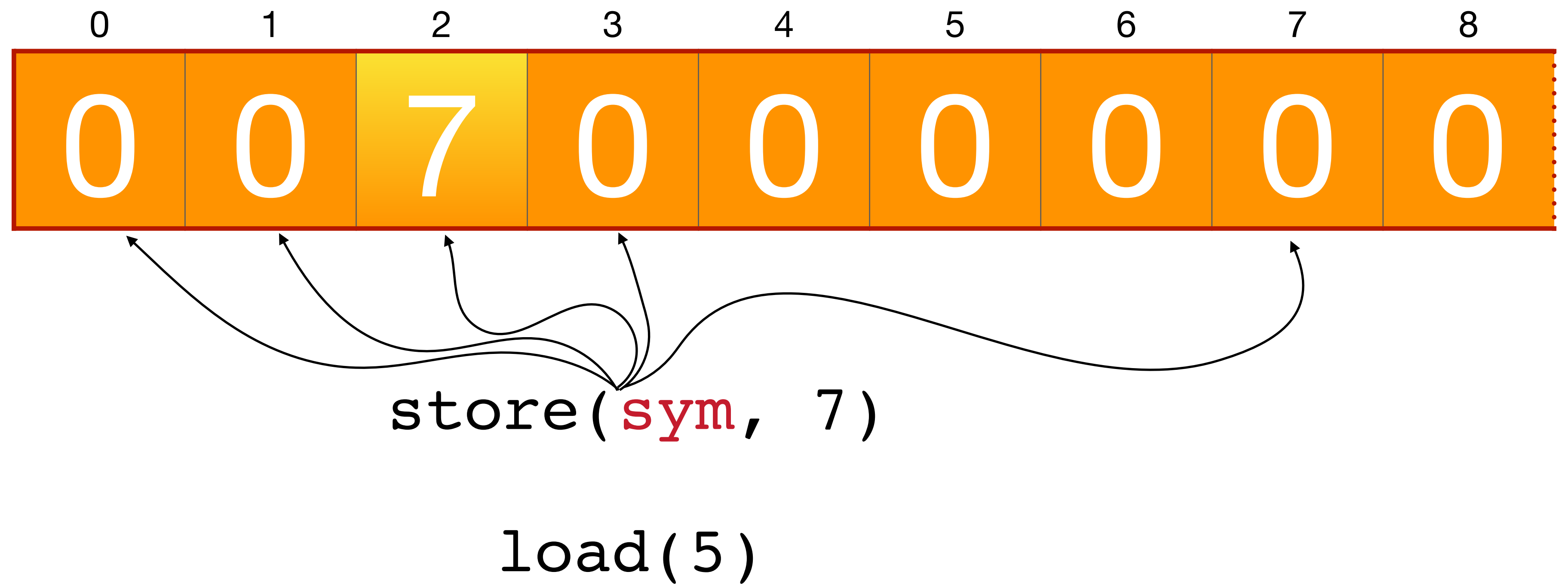
Handling Symbolics



Handling Symbolics



Handling Symbolics



Fine-grained Memory-Object Representation

**Insight I:
Differences are small and
common parts large**

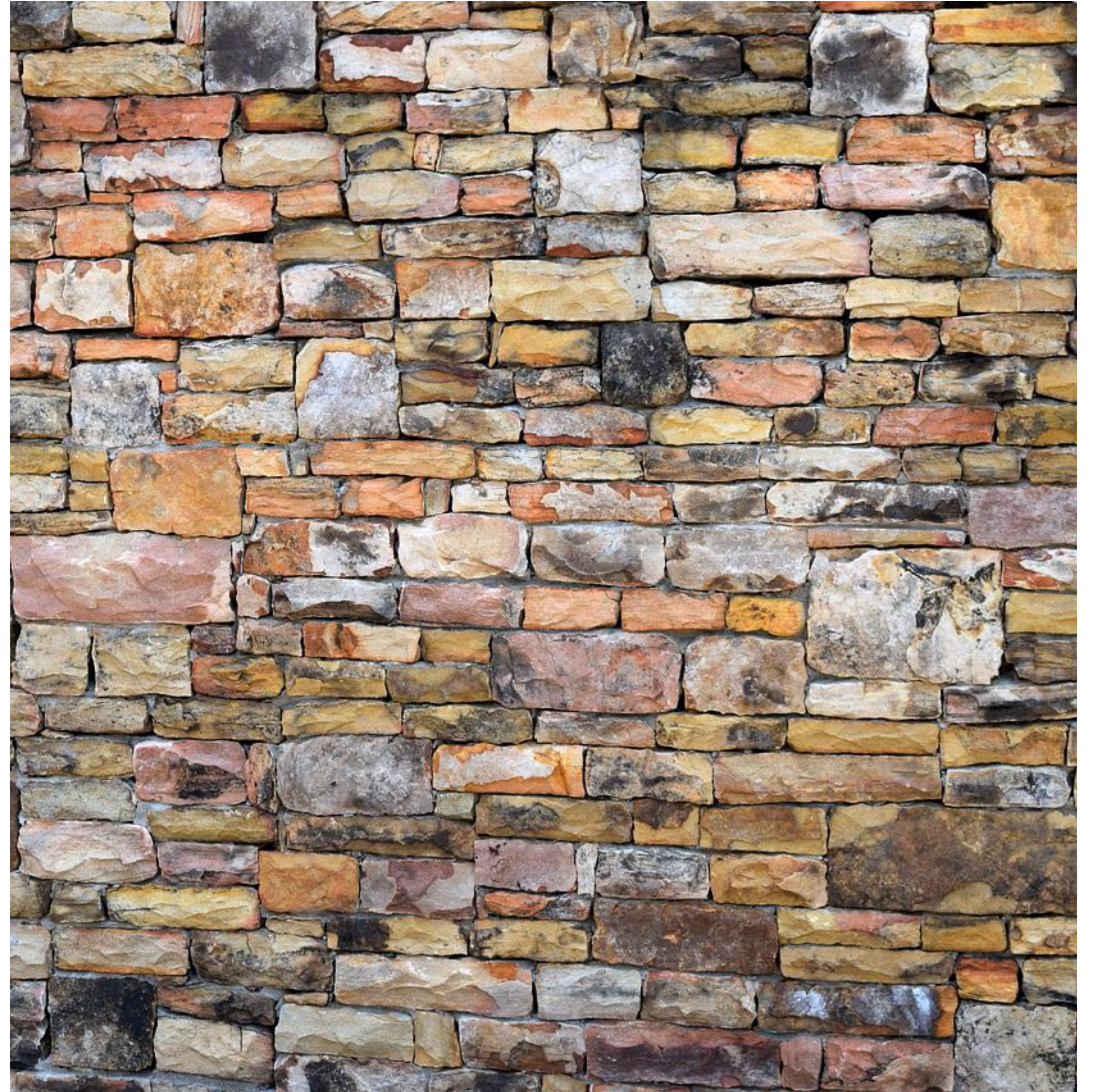
Insight II:
Changes are (often) **local** and
of similar type

←-----→ Space

⌚ Time

Basic Building Blocks

Everything is a Layer

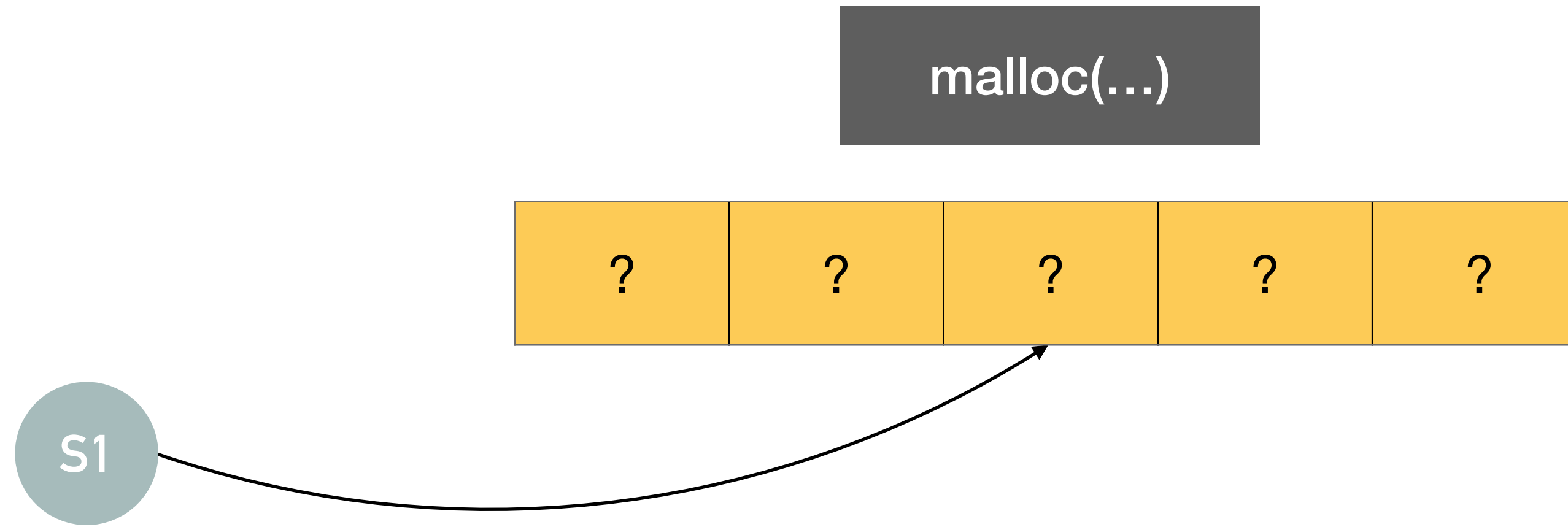


Example Scenario

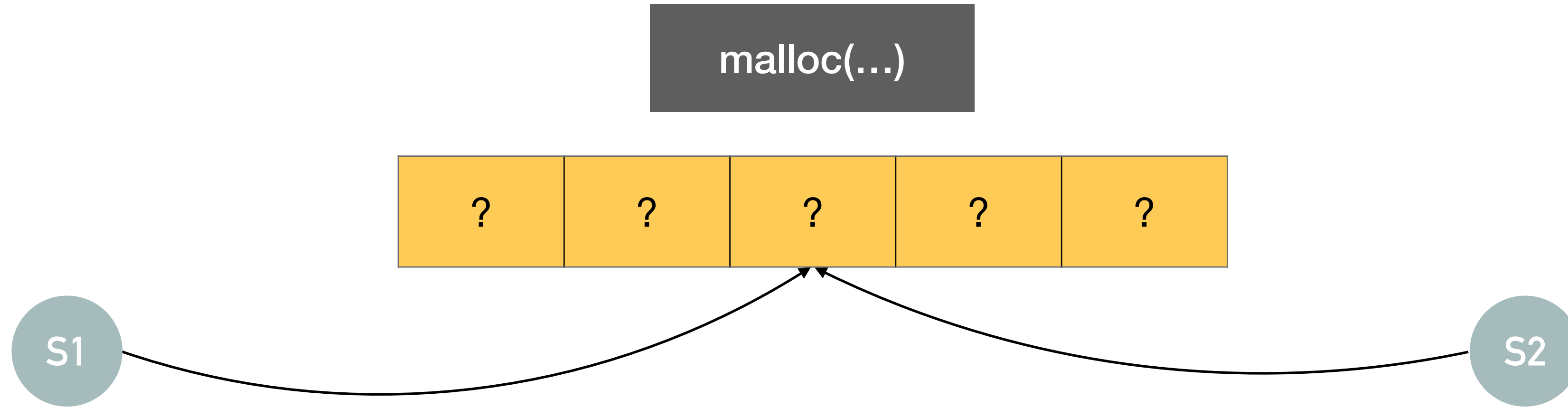
malloc(...)

S1

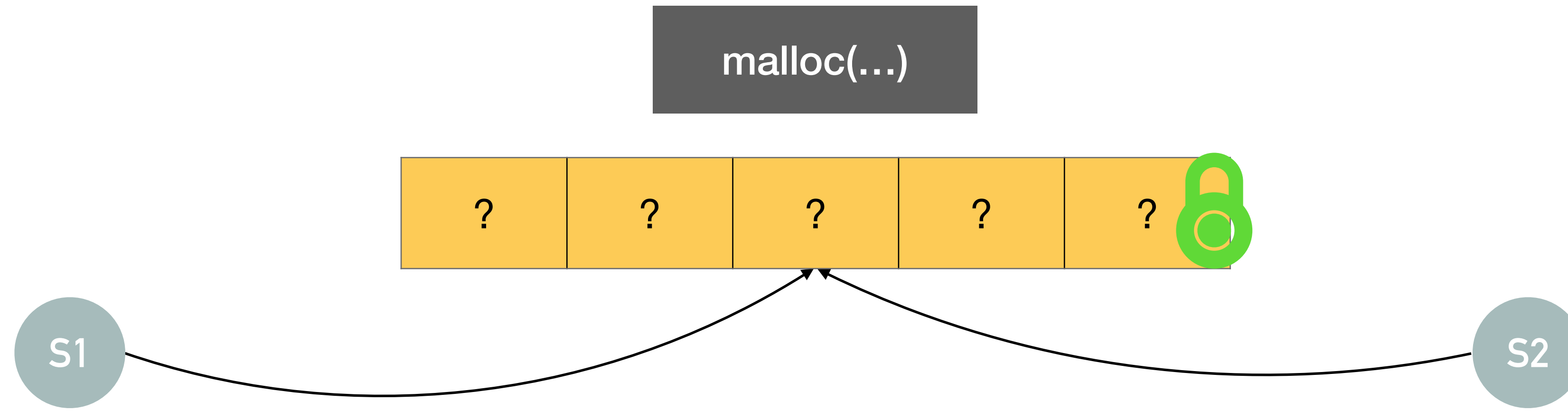
Example Scenario



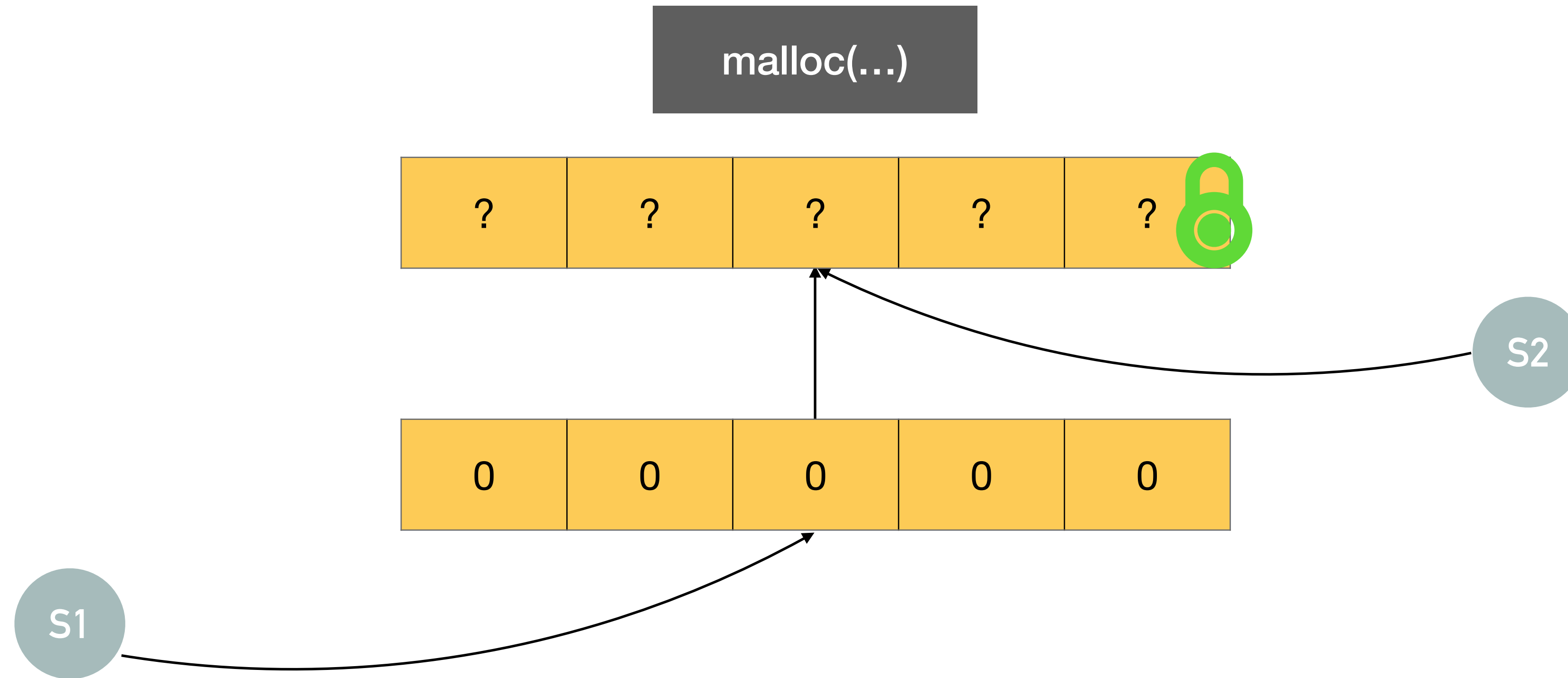
Example Scenario



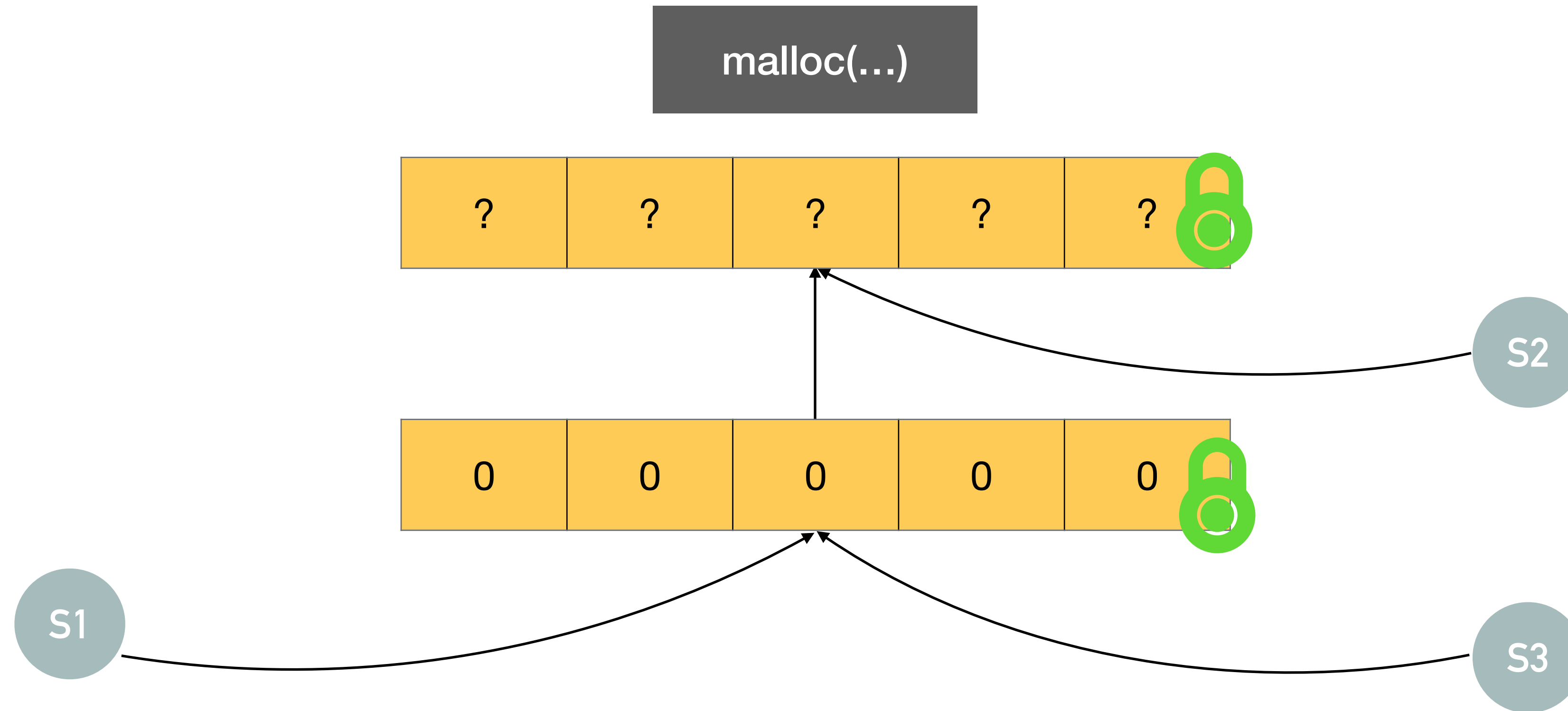
Example Scenario



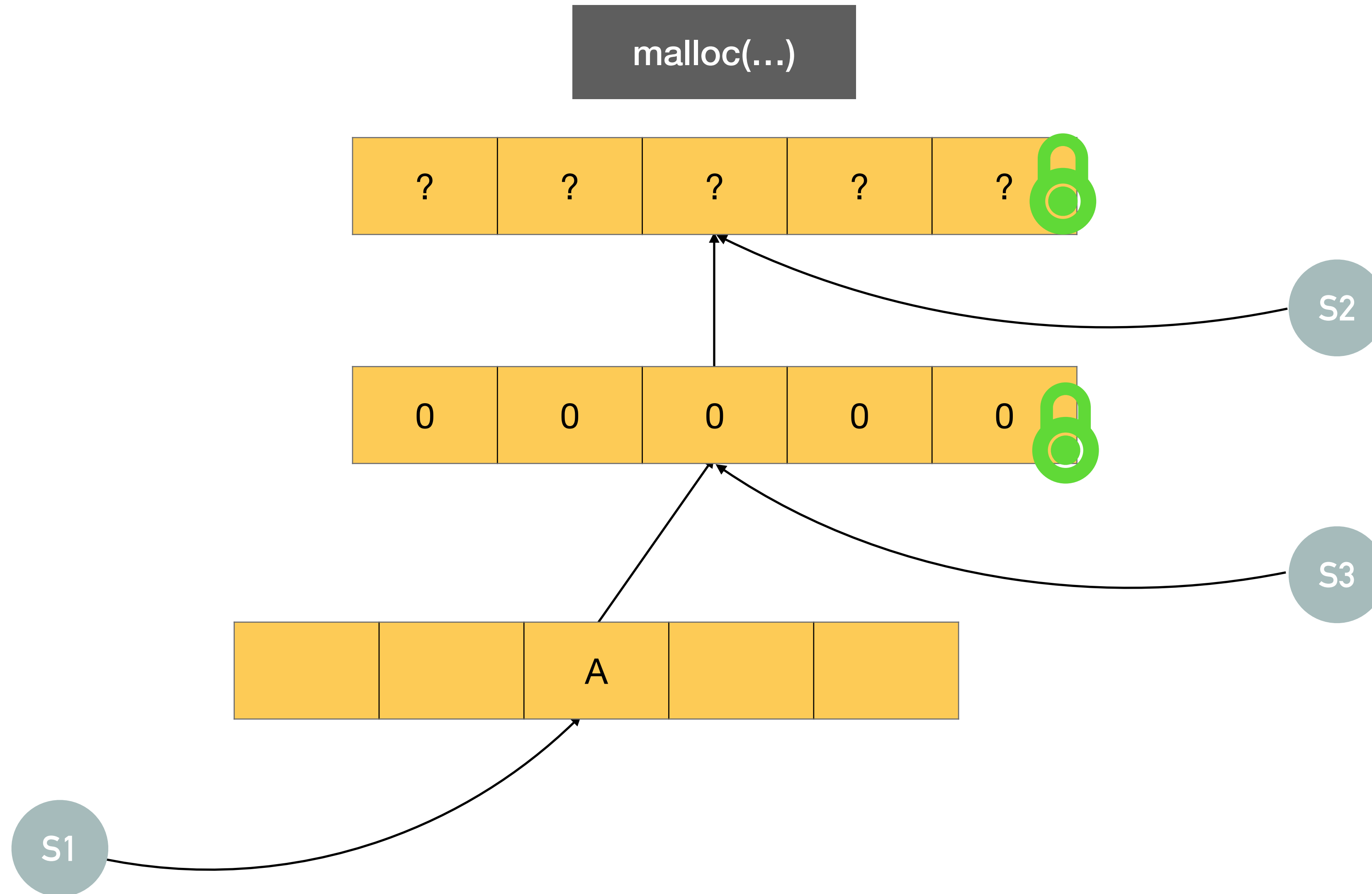
Example Scenario



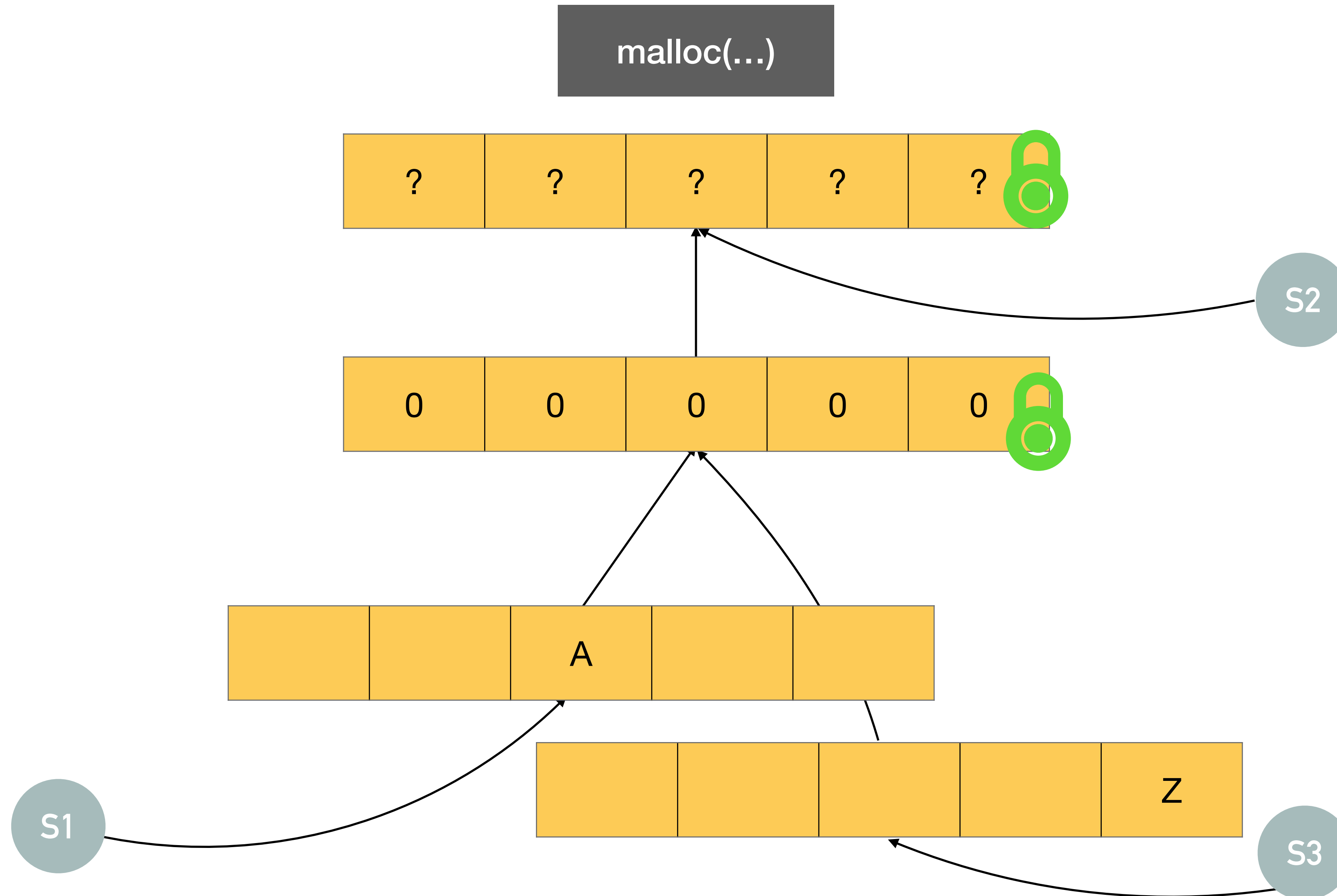
Example Scenario



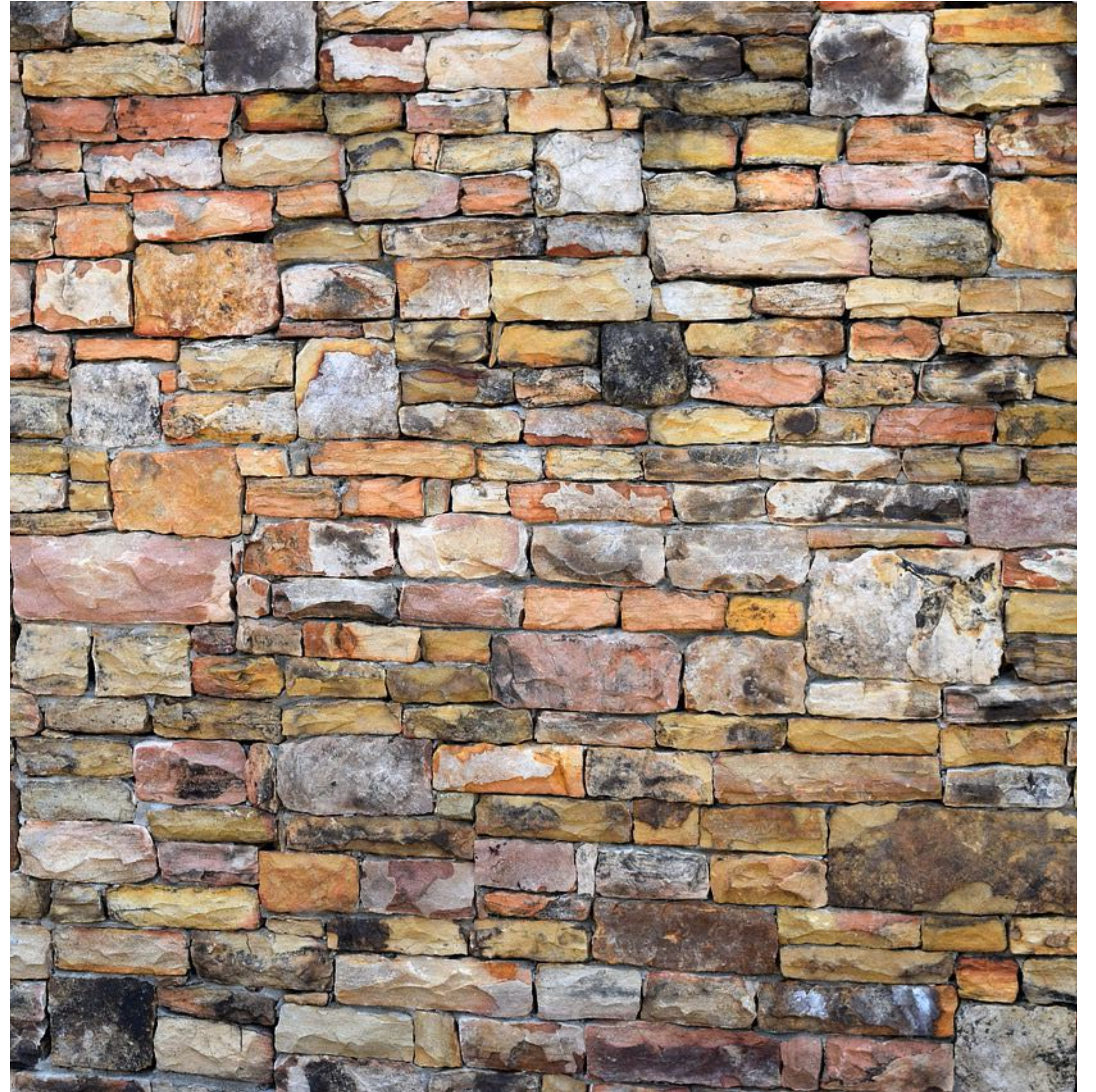
Example Scenario



Example Scenario

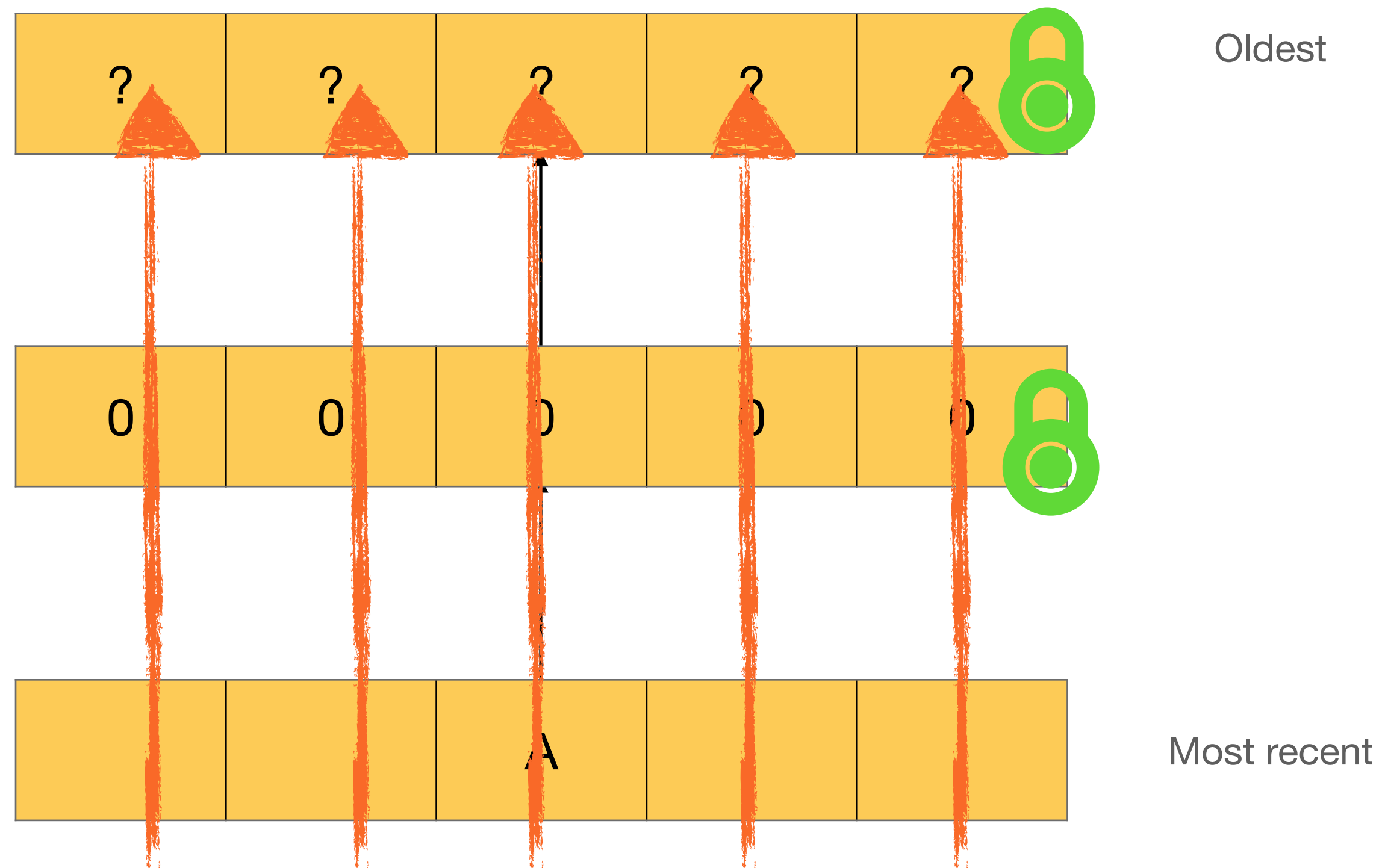


Optimisations



Index-based Access

malloc(...)

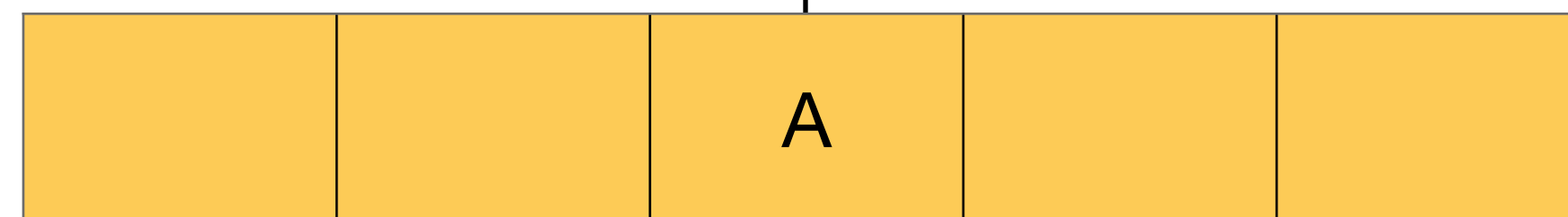
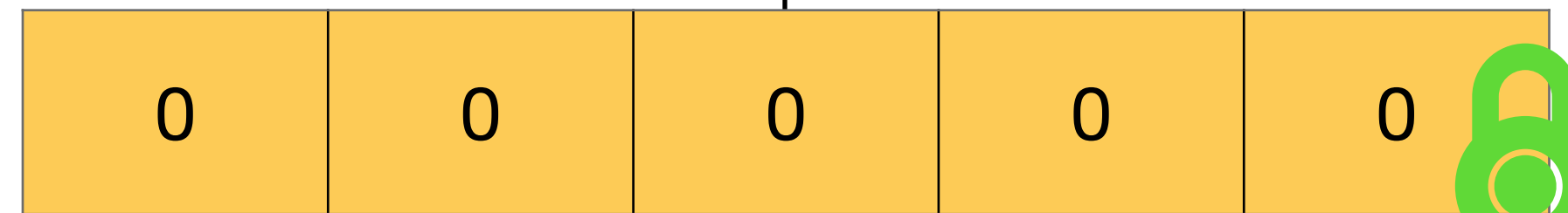


Index-based Access

malloc(...)



Oldest



Most recent

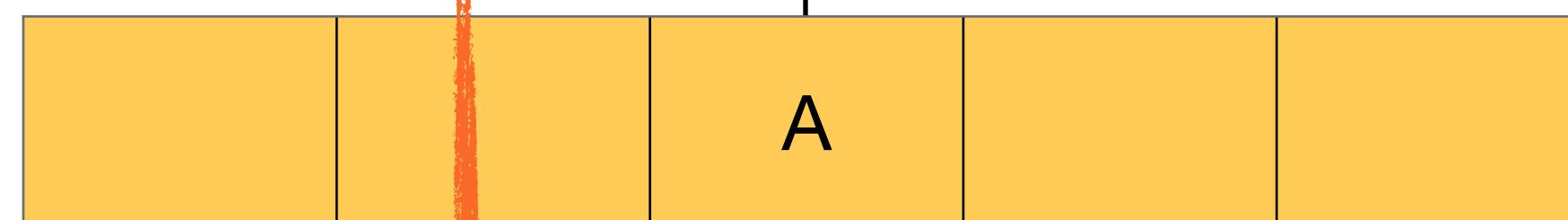
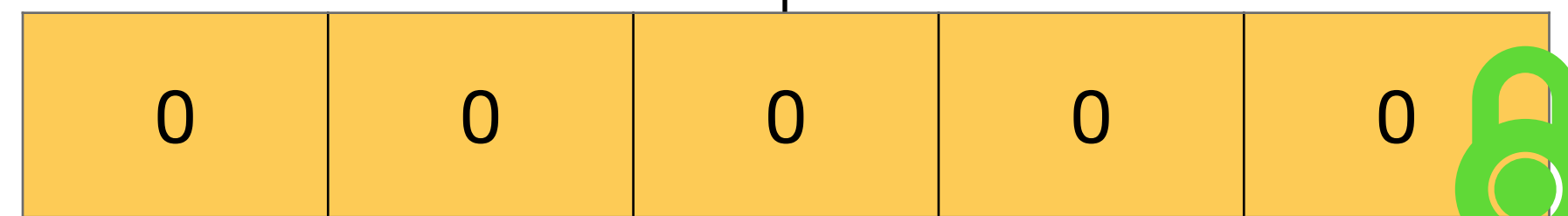
load(2) -> A

Index-based Access

malloc(...)



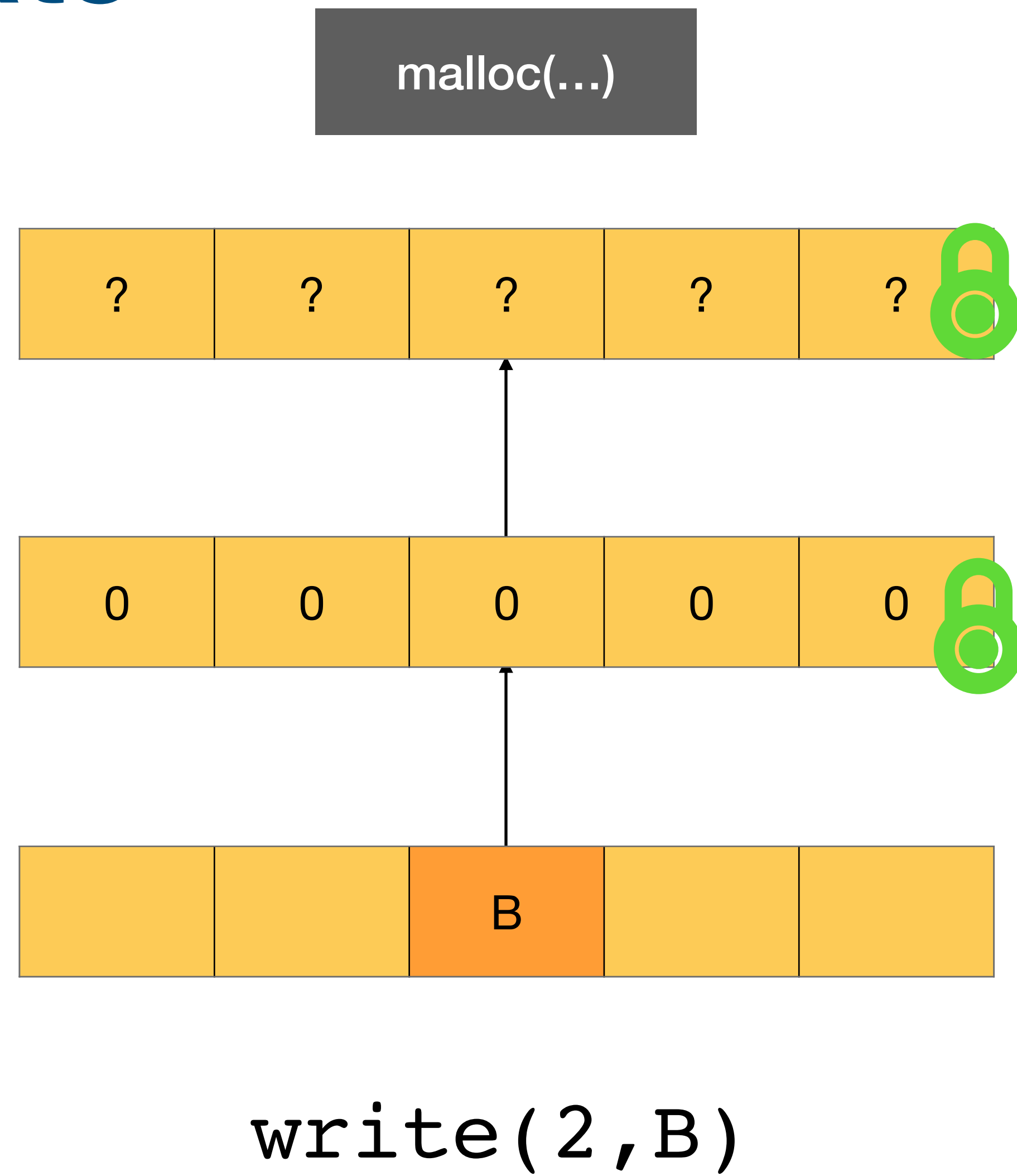
Oldest



Most recent

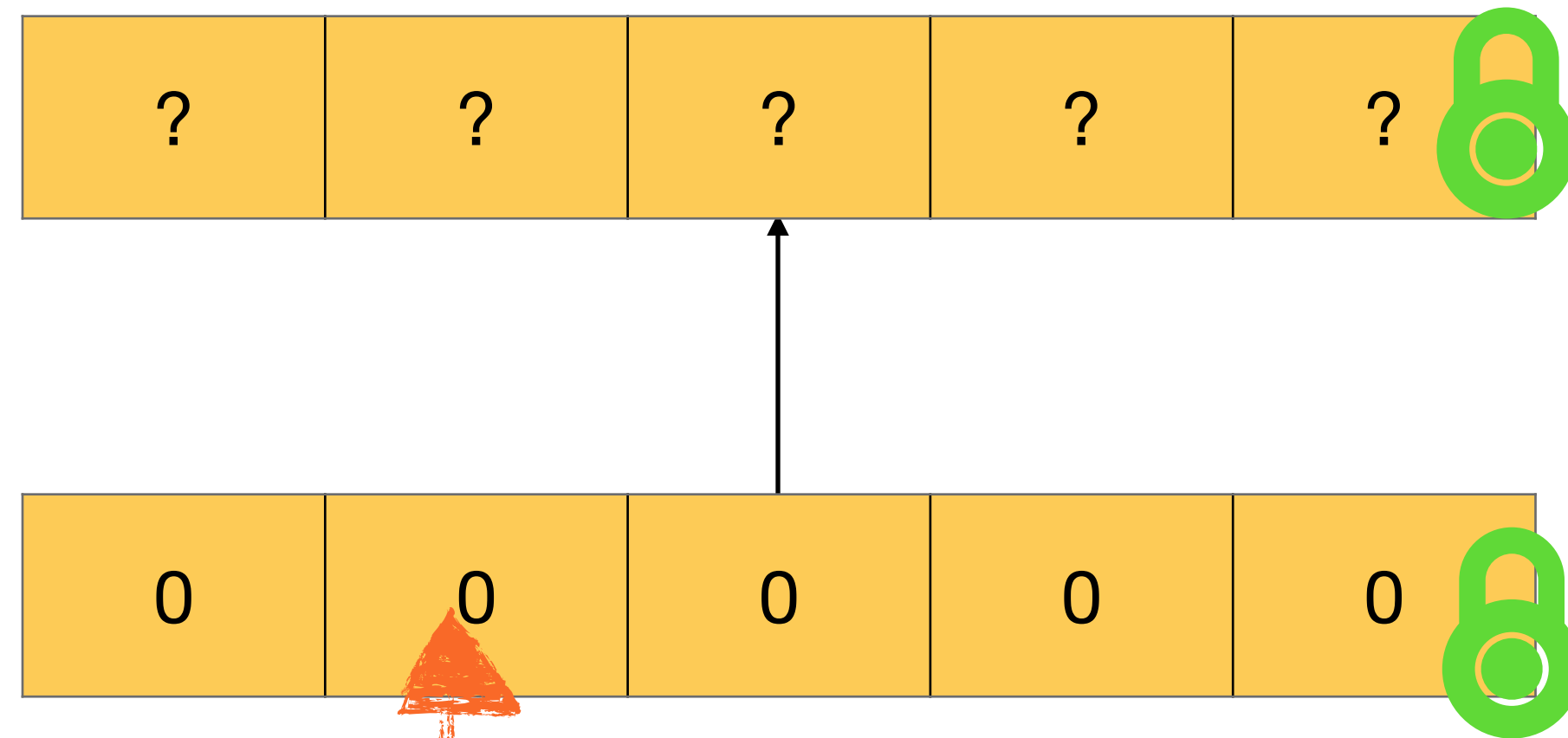
load(1) -> 0

In-place update



Conditional update

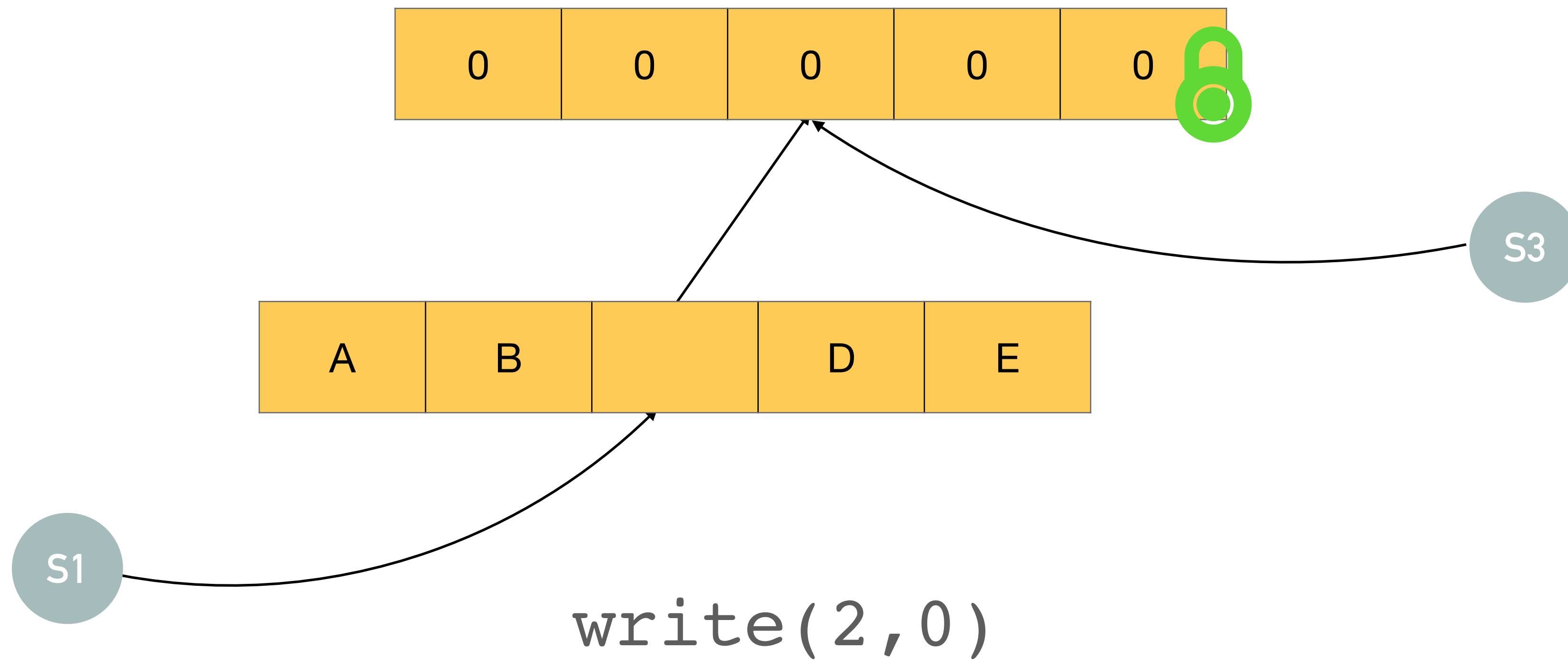
`malloc(...)`



`write(1,0)`

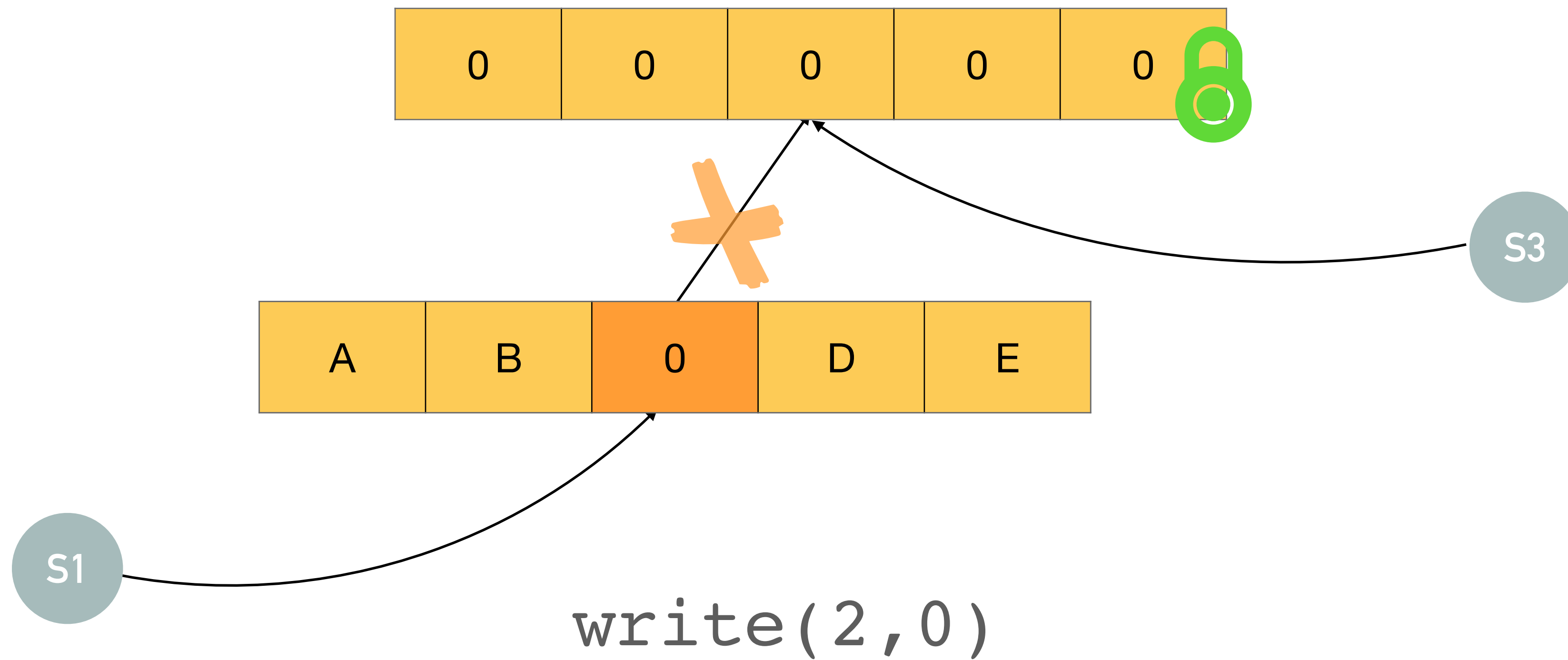
Layer Invalidation

`malloc(...)`



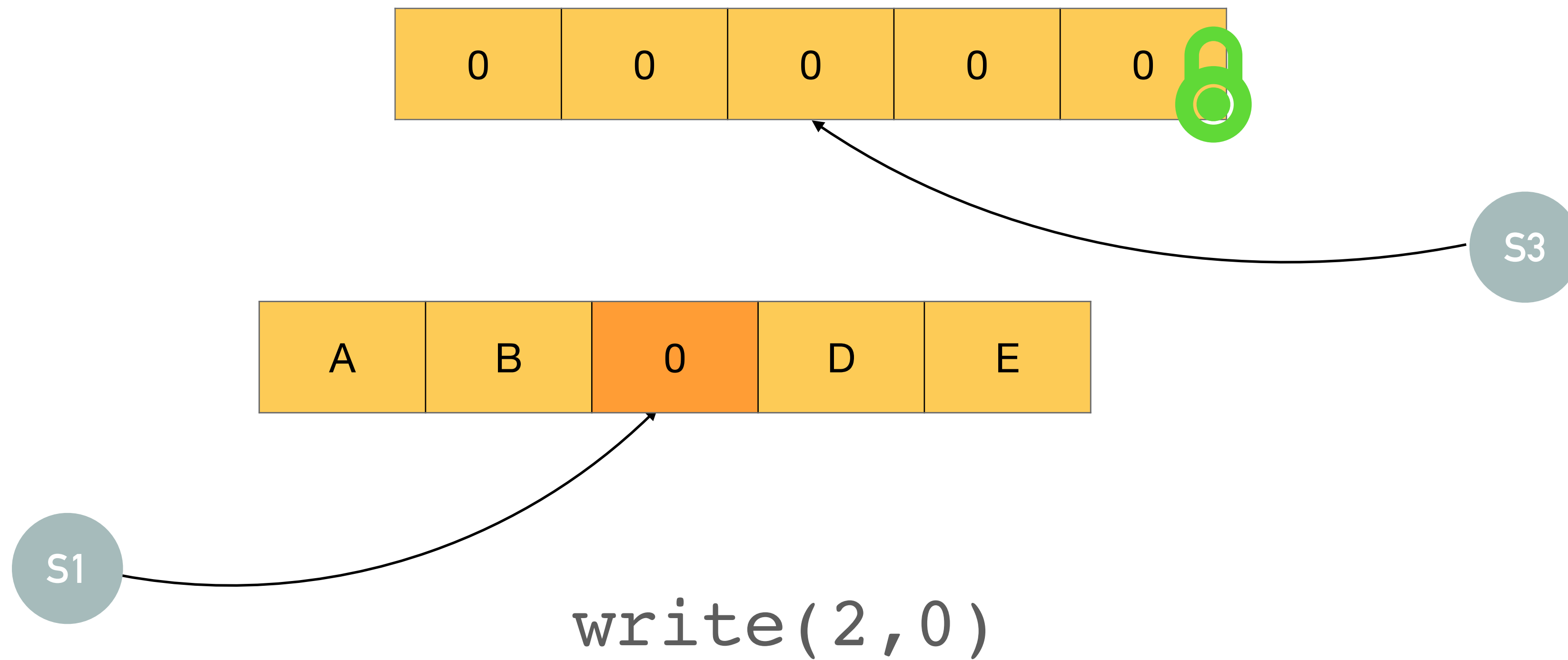
Layer Invalidation

`malloc(...)`



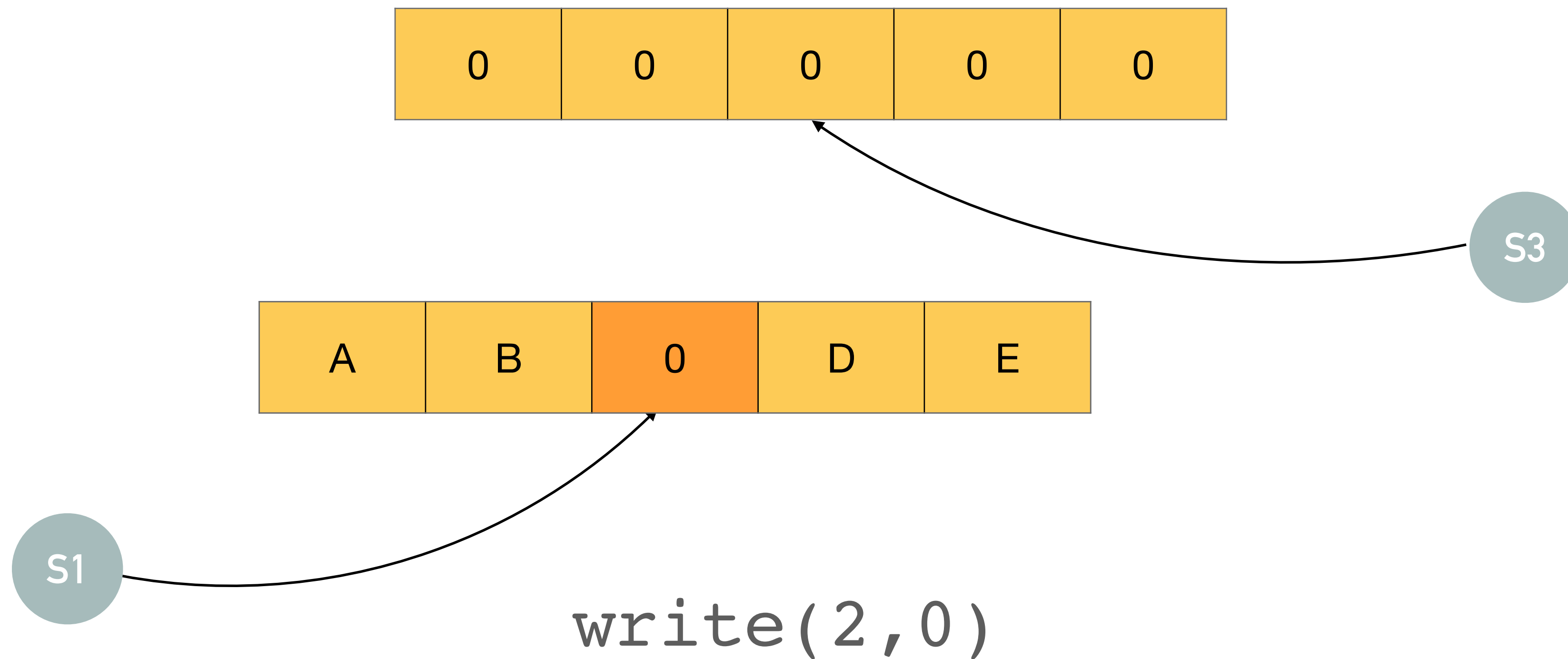
Layer Invalidation

`malloc(...)`

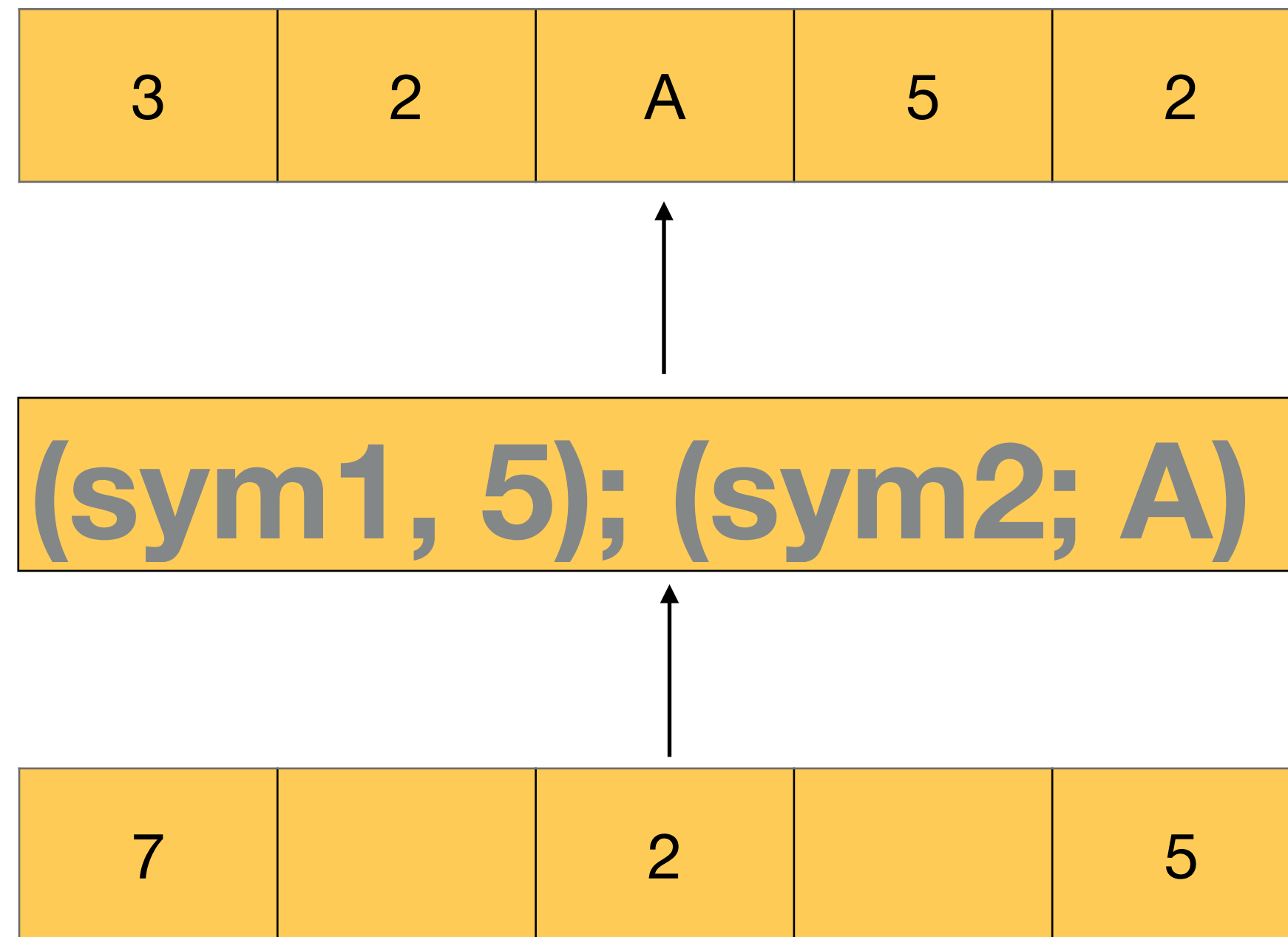


Layer Invalidation

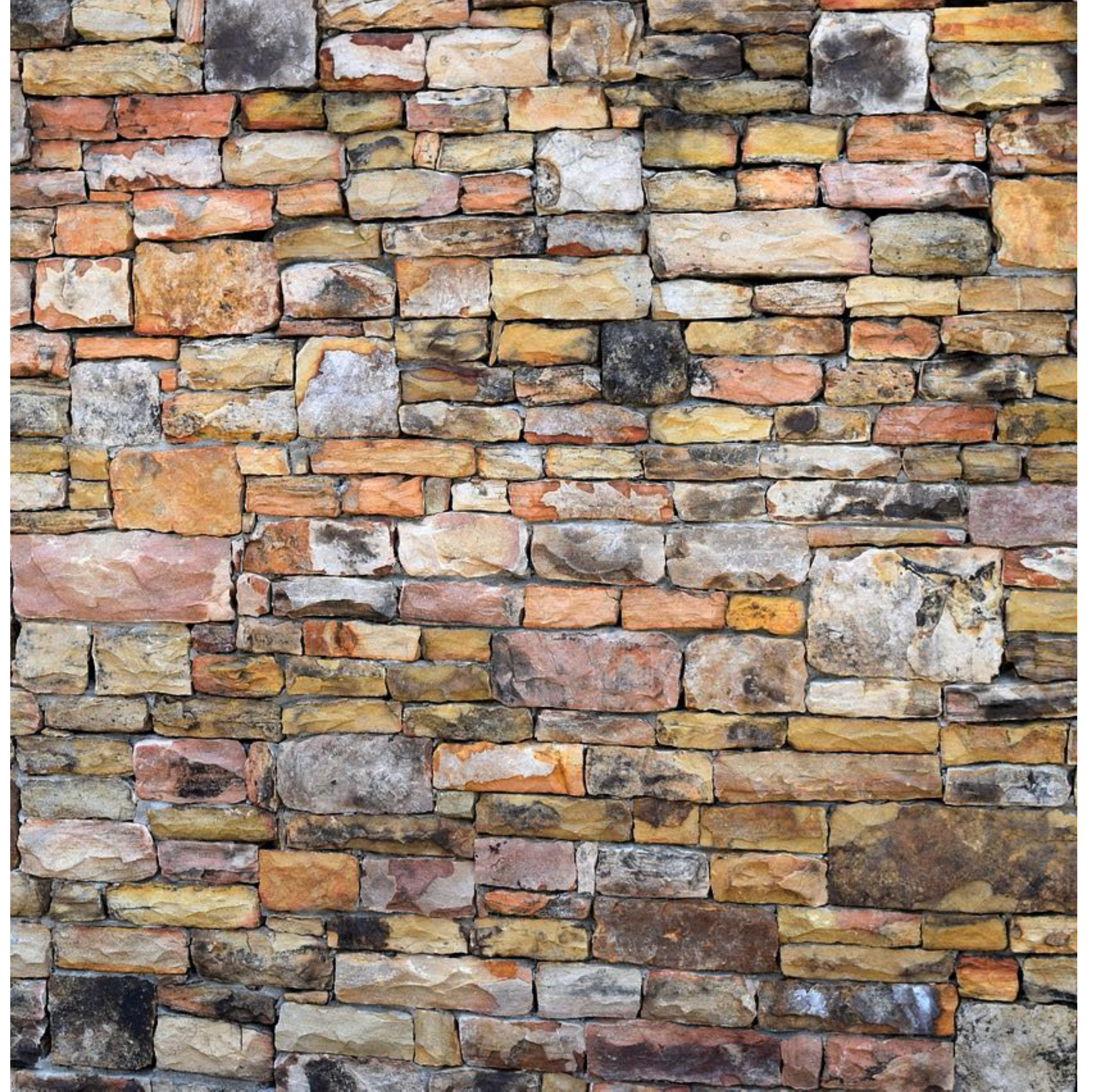
`malloc(...)`



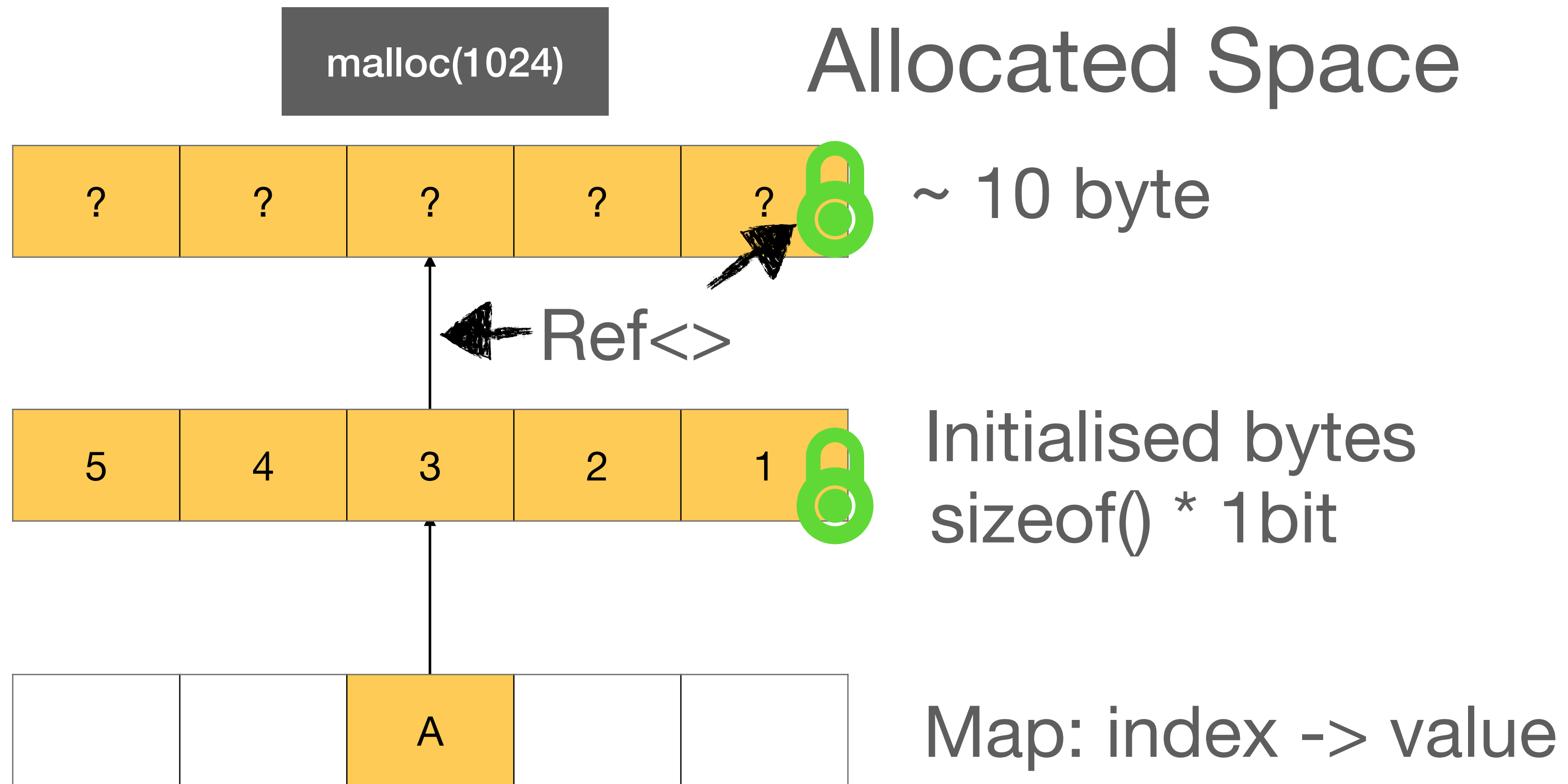
Handling Symbolic Indices



Implementation



Layer Types

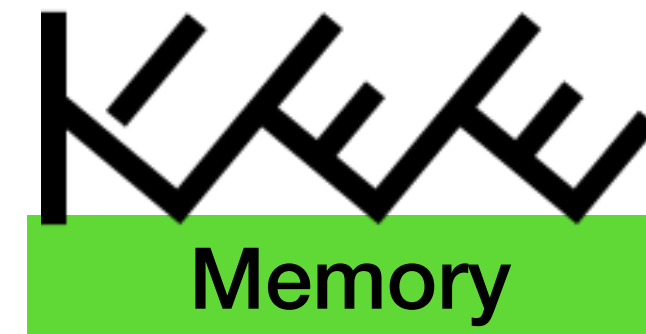


Evaluation

Benchmarks



vs.

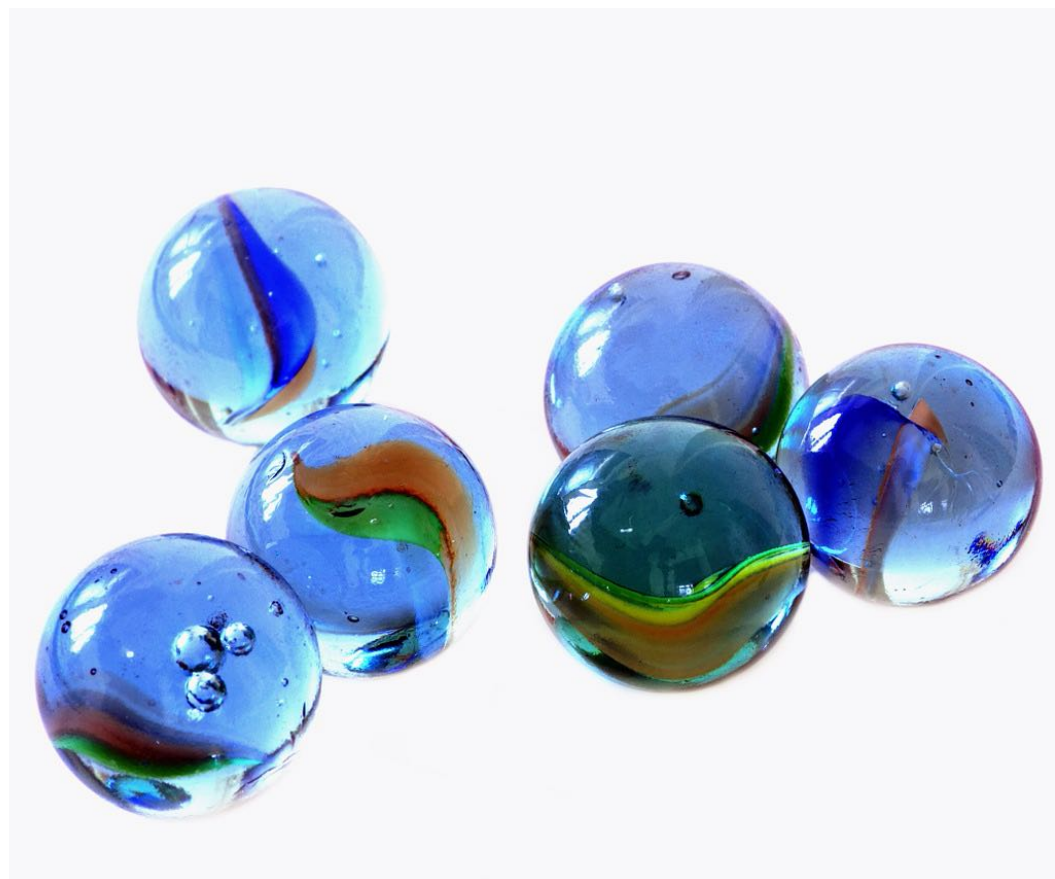


GNU Coreutils

Depth-First

Search
Strategies

Breadth-First



Random +
Target Uncovered

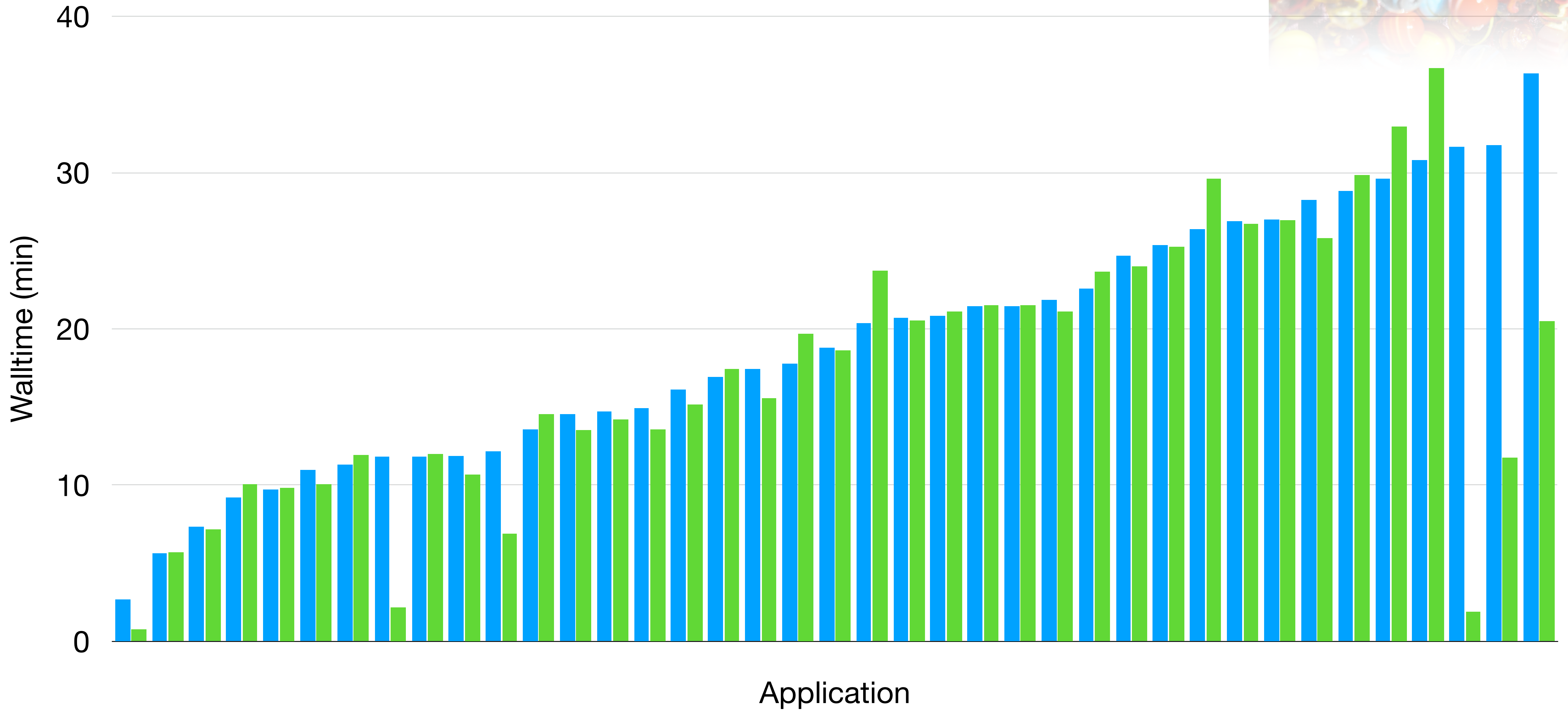


RQ1: Changes in Execution Time

Walltime - Breadth First Search

■ KLEE

■ Memory

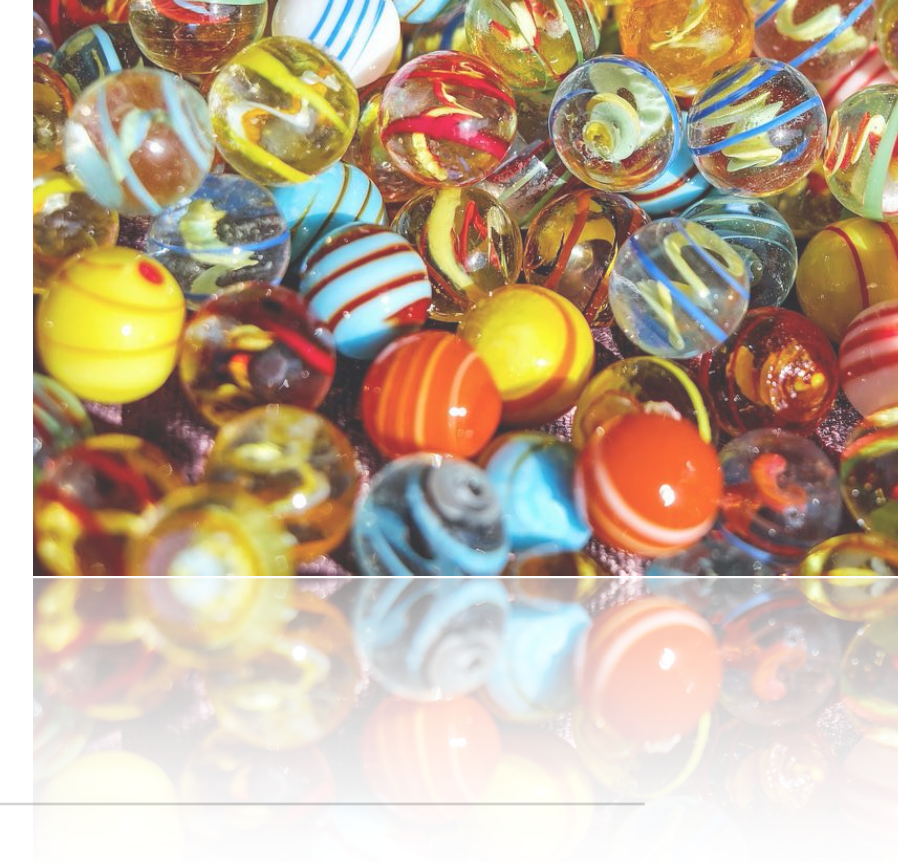
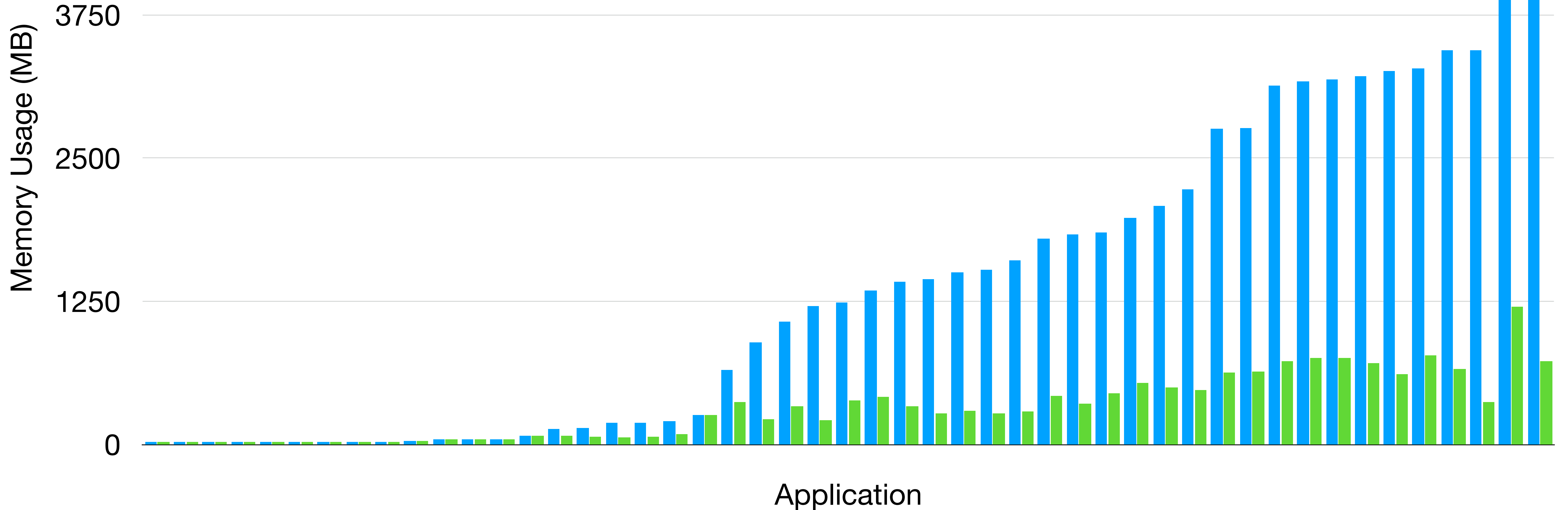


RQ2: Changes in Memory Consumption

Memory Usage - Breadth First Search

KLEE Memory

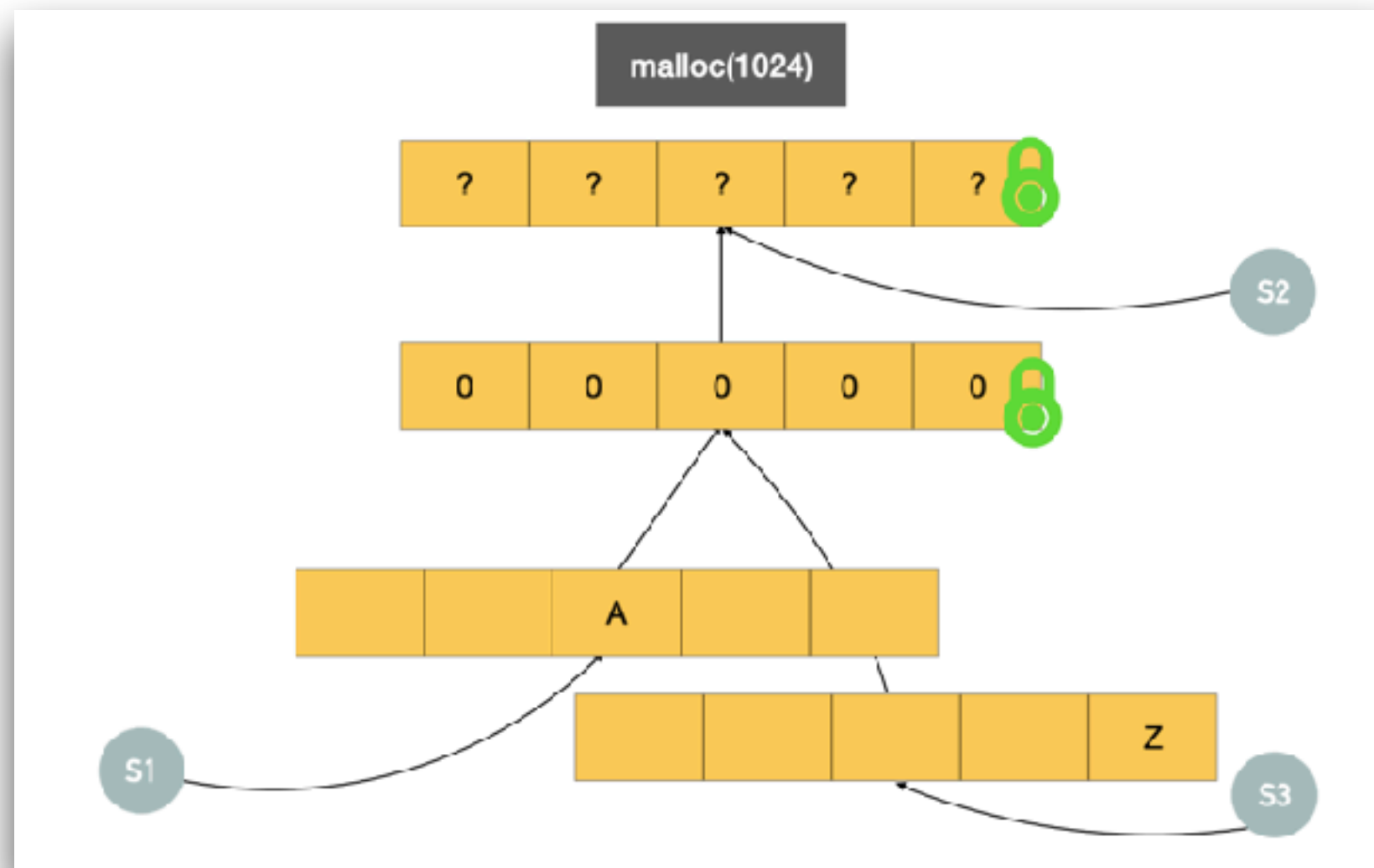
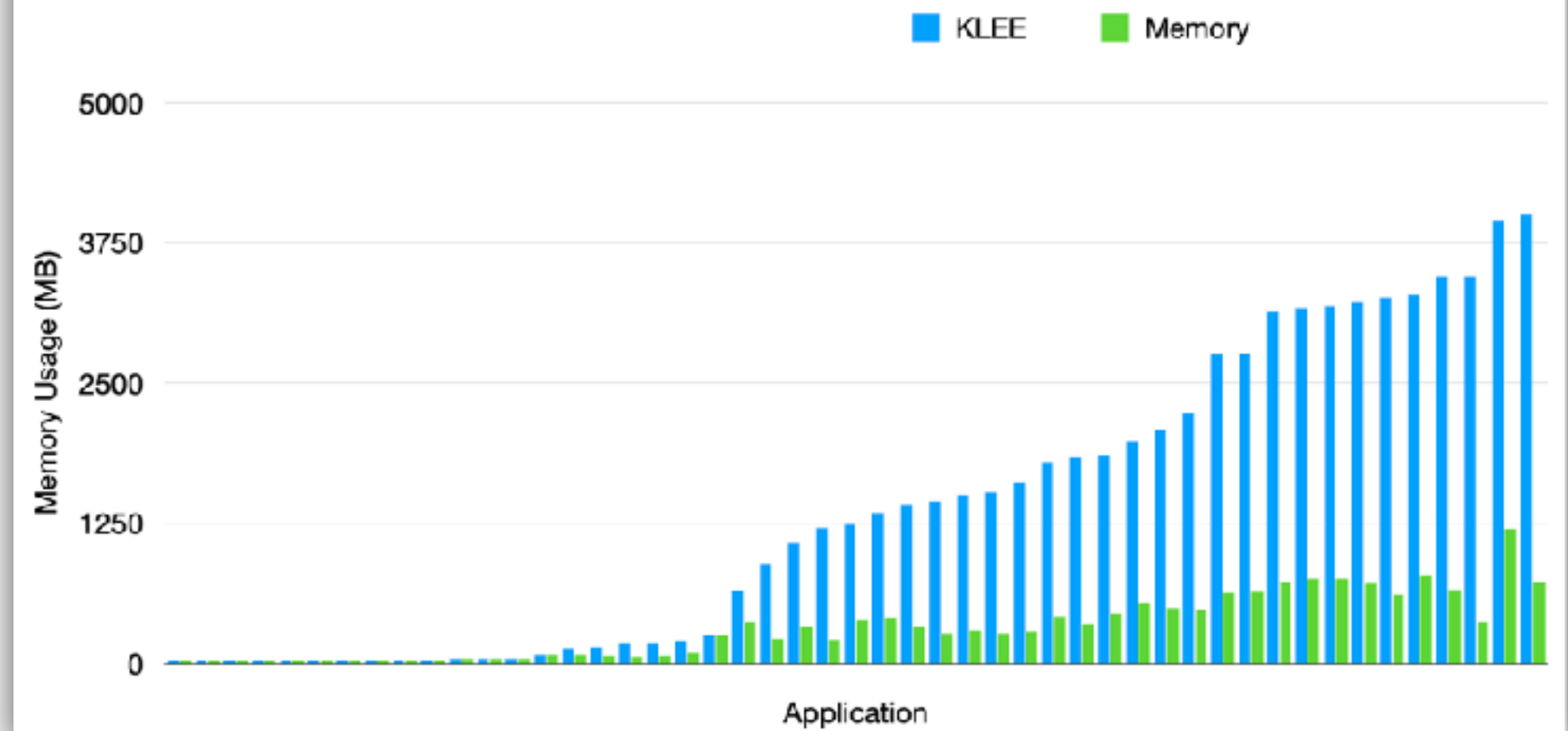
5000





Summary

Memory Usage - Breadth First Search



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