

TracerX-Pruning Dynamic Symbolic Execution with Weakest Precondition Interpolation

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From KLEE To TracerX

- DFS Forward Symbolic Execution to find feasible paths (Similar to KLEE)
- Intermediate execution states preserved (Unlike KLEE)
- Path interpolants are generated for each path during backward tracking
- Tree interpolants are generated as conjunction of path interpolants
- Tree interpolants then used for subsumption at similar program points

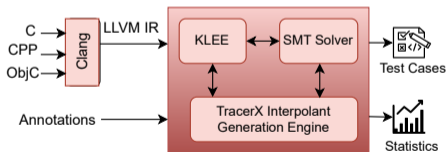
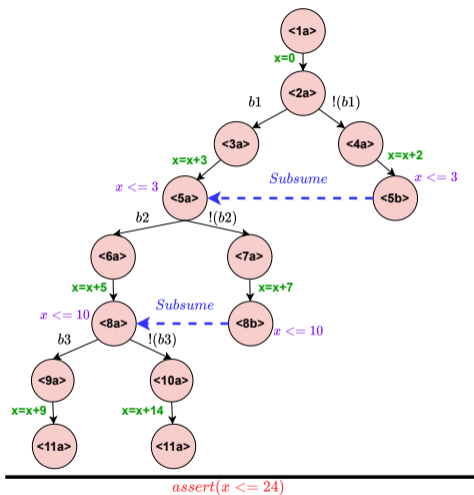


Figure: TracerX Framework



Figure: Pruning of subtree

Symbolic Execution Tree with Interpolation



```
x = 0;  
if (b1) x += 3 else x += 2  
if (b2) x += 5 else x += 7  
if (b3) x += 9 else x += 14  
assert(x <= 24)
```

- **Without interpolation:** The full tree is traversed.
- **With interpolation:**
 - 1 $\langle 8b \rangle$ context contains $x = 10$. It is subsumed with the tree interpolant from $\langle 8a \rangle$: $x \leq 10$.
 - 2 $\langle 5b \rangle$ context contains $x = 2$. Subsumed with the tree interpolant from $\langle 5a \rangle$: $x \leq 3$.
 - 3 Big subtree traversal is avoided.

Interpolation: Weakest Precondition

- **Ideal interpolant** is the weakest precondition (WP) of the target. Unfortunately, WP is **intractable** to compute.
- For example, Assume $(b1 \wedge \neg b2 \wedge \neg b3)$ is **UNSAT**.
WP before first “if-statement” is: $b1 \rightarrow (\neg b2 \wedge b3 \wedge x \leq 7) \vee (b2 \wedge x \leq 4)$
 $\neg b1 \rightarrow x < 3$
- Essentially, WP is **exponentially disjunctive**
- Challenge is to obtain a conjunctive approximation

A Path is a sequence of assignment and assume instructions:

- 1 Interpolant of **Assignment** instruction:
 - $WP(inst, \omega) = \dots$ inverse transition of $inst$ over ω
 - e.g. $\omega : x \leq 15$ and $inst : x = z + 2$, then $WP(inst, \omega) : z \leq 13$
- 2 Interpolant of **Assume** instruction (C is incoming Context): $\{C\}$ assume(B) $\{\omega\}$
 - WP Approximation: find \bar{C} to replace C
 - ABDUCTION PROBLEM !!!

Approximation of Weakest Precondition

This algorithm is the **heart of TracerX**:

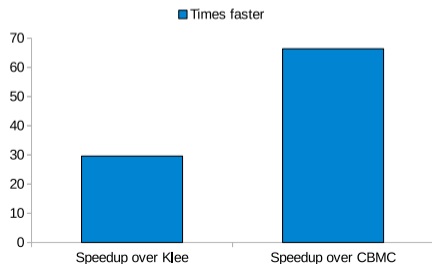
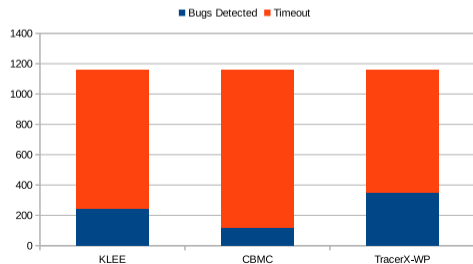
- 1 We compute finest partition so that $\text{var}(C_i) * \text{var}(C_j)$ s.t. $i \neq j$:
 $\{C_1 * C_2 * C_3 * \dots * C_n\}$ $\text{assume}(B)$ $\{\omega_1 * \omega_2 * \omega_3 * \dots * \omega_m\}$ (* is as in separation logic).
- 2 Bunch C_i into three:

- **Target independent:** The C_i which are separate from B and ω .
Action: Replace C_i with *true*, i.e. remove C_i .
- **Guard independent:** Consider $C_{gi} \equiv C_i$ s.t. $C_i * B$; and, $\omega_{gi} \equiv \omega_j$ s.t. $B * \omega_j$.
Action: Replace C_{gi} by ω_{gi} .
- **Remainder of the C_i :** We do not capture exact WP for this group.
e.g. $\{z == 5\}$ $\text{assume}(x > z - 2)$ $\{x > 0\}$ (Here, $z > 2$ is the WP)
Action: No change to C_i , i.e. keep C_i .

Experimental Results

Data set: All C-programs from RERS-2012 Challenge [6].

- **Total targets:** 1159
- All three systems **KLEE** [1], **CBMC** [5] and **TracerX-WP** [4] are run for 3600 seconds
- ① TracerX-WP able to detect **348** targets, while KLEE and CBMC are detected **245** and **117** targets respectively.
- ② TracerX-WP is **29.59x** faster than KLEE and **66.37x** faster than CBMC



- 1 **Website:** <https://tracer-x.github.io/>
- 2 **Github:** <https://github.com/tracer-x/>
- 3 **TracerX: Dynamic Symbolic Execution with Interpolation**
J. Jaffar, R. Maghareh, S. Godbole, X.L. Ha, 2020
<https://arxiv.org/abs/2012.00556>
- 4 **TracerX: Dynamic Symbolic Execution with Interpolation (competition contribution)** J. Jaffar, R. Maghareh, S. Godbole, X.L. Ha,
- 5 **Toward Optimal MC/DC Test Case Generation**
S. Godbole, J. Jaffar, R. Maghareh, A. Dutta, *ISSTA 2021*
- 6 **TracerX: Pruning Dynamic Symbolic Execution with Deletion and Weakest Precondition Interpolation (competition contribution)**
A. Dutta, R. Maghareh, J. Jaffar, S. Godbole, X. L. Yu, *FASE 2024*

- [1] C. Cadar et al. Klee: Unassisted and automatic generation of high-coverage tests for complex systems programs. In: OSDI, 2008.
- [2] J. Jaffar et al. TRACER: A symbolic execution tool for verification. In: CAV, 2012.
- [3] J. Jaffar et al. TracerX: Dynamic symbolic execution with interpolation (competition contribution) . In: FASE, 2020.
- [4] A. Dutta et al. TracerX: Pruning Dynamic Symbolic Execution with Deletion and Weakest Precondition Interpolation (competition contribution). In: FASE, 2024.
- [5] D. Kroening D et al. CBMC-C Bounded Model Checker. In: TACAS 2014.
- [6] <http://rers-challenge.org/>